

# Note about the linear complexity of new generalized cyclotomic binary sequences of period $2p^n$

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*Abstract:* This paper examines the linear complexity of new generalized cyclotomic binary sequences of period  $2p^n$  recently proposed by Yi Ouang et al. (arXiv:1808.08019v1 [cs.IT] 24 Aug 2018). We generalize results obtained by them and discuss author's conjecture of this paper.

*Key-Words:* Binary sequences, linear complexity, cyclotomy

## 1 Introduction

The cyclotomic classes and the generalized cyclotomic classes are often used for design sequences with high linear complexity, which is an important characteristic of sequence for the cryptography applications [2]. Recently, new generalized cyclotomic classes were presented in [8]. The linear complexity of new generalized cyclotomic binary sequences with period  $p^n$  was studied in [9, 4, 7]. A new family of binary sequences with period  $2p^n$  based on the generalized cyclotomic classes from [8] was presented in [6]. Yi Ouang et al. examined the linear complexity of these sequences for  $f = 2^r$ , where  $p = 1 + ef$  and  $r$  is a positive integer. They offered new studying method of the linear complexity of these sequences. Their method based on ideas from [4].

In this paper we show that for study of the linear complexity of new sequence family from [6] we can use only old the method from [4]. Furthermore, it will be enough for obtaining more generalized results than in [6] and for the proof and the correction of the conjecture of the authors of this paper. Here we keep the notation and the structure of [4], i.e., in Sect. 2 we introduce some basics and recall the definition of a generalized cyclotomic sequence and the conjecture from [6]. Section 3 is dedicated to the study of the linear complexity of this family of cyclotomic sequences. Section 4 concludes the work in this paper.

We will study the linear complexity of new sequence family from [6] when  $p$  is not a Wieferich prime, i.e.  $2^{p-1} \not\equiv 1 \pmod{p^2}$ . It was shown that there are only two such primes, 1093 and 3511, up to  $6 \times 10^{17}$  [1, 3].

## 2 Preliminaries

Throughout this paper, we will denote by  $\mathbb{Z}_N$  the ring of integers modulo  $N$  for a positive integer  $N$ , and by  $\mathbb{Z}_N^*$  the multiplicative group of  $\mathbb{Z}_N$ .

First of all we will recall some basics of the linear complexity of a periodic sequence and introduce the generalized cyclotomic sequences proposed in [6].

### 2.1 Linear Complexity

Let  $s^\infty = (s_0, s_1, s_2, \dots)$  be a binary sequence of period  $N$  and  $S(x) = s_0 + s_1x + \dots + s_{N-1}x^{N-1}$ . It is well known (see, for instance, [2, Page 171]) that the linear complexity of  $s^\infty$  is given by

$$L(s^\infty) = N - \deg \left( \gcd(x^N - 1, S(x)) \right).$$

So, if  $N = 2p^n$  then we see that

$$L(s^\infty) = 2p^n - \deg \left( \gcd((x^{p^n} - 1)^2, S(x)) \right).$$

Thus, if  $\alpha_n$  is a primitive root of order  $p^n$  of unity in the extension of the field  $\mathbb{F}_2$  (the finite field of two elements) then in order to find the linear complexity of a sequence it is sufficient to find the zeros of  $S(x)$  in the set  $\{\alpha_n^i, i = 0, 1, \dots, p^n - 1\}$  and determine their multiplicity.

### 2.2 New Generalized Cyclotomic Sequences Length $2p^n$

Let  $p$  be an odd prime and  $p = ef + 1$ , where  $e, f$  are positive integers. Let  $g$  be a primitive root modulo

$p^n$ . It is well known [5] that an odd number from  $g$  or  $g + p^n$  is also a primitive root modulo  $2p^j$  for each integer  $j \geq 1$ . Hence, we can assume that  $g$  is an odd number. Further, the order of  $g$  modulo  $2p^j$  is equal to  $\varphi(2p^j) = p^{j-1}(p - 1)$ , where  $\varphi(\cdot)$  is the Euler's totient function. Below we recall the definitions of generalized cyclotomic classes introduced in [8] and [6].

Let  $n$  be a positive integer. For  $j = 1, 2, \dots, n$ , denote  $d_j = p^{j-1}f$  and define

$$D_0^{(p^j)} = \left\{ g^{t \cdot d_j} \pmod{p^j} \mid 0 \leq t < e \right\}, \text{ and}$$

$$D_i^{(p^j)} = g^i D_0^{(p^j)} = \left\{ g^i x \pmod{p^j} : x \in D_0^{(p^j)} \right\},$$

$$1 \leq i < d_j,$$

$$D_0^{(2p^j)} = \left\{ g^{t \cdot d_j} \pmod{2p^j} \mid 0 \leq t < e \right\}, \text{ and}$$

$$D_i^{(2p^j)} = g^i D_0^{(2p^j)} = \left\{ g^i x \pmod{2p^j} : x \in D_0^{(2p^j)} \right\},$$

$$1 \leq i < d_j. \quad (1)$$

The cosets  $D_i^{(p^j)}$ ,  $i = 0, 1, \dots, d_j - 1$ , are called *generalized cyclotomic classes* of order  $d_j$  with respect to  $p^j$ . It was shown in [8] that  $\{D_0^{(p^j)}, D_1^{(p^j)}, \dots, D_{d_j-1}^{(p^j)}\}$  forms a partition of  $\mathbb{Z}_{p^j}^*$  for each integer  $j \geq 1$  and for an integer  $m \geq 1$ . Also  $\{D_0^{(2p^j)}, D_1^{(2p^j)}, \dots, D_{d_j-1}^{(2p^j)}\}$  forms a partition of  $\mathbb{Z}_{2p^j}^*$  for each integer  $j \geq 1$  and for an integer  $m \geq 1$ .

Let  $f$  be a positive even integer and  $b$  an integer with  $0 \leq b < p^{n-1}f$ . Define four sets

$$C_0^{(2p^n)} = \bigcup_{j=1}^n \bigcup_{i=d_j/2}^{d_j-1} p^{n-j} \left( D_{(i+b)}^{(2p^j)} \pmod{d_j} \right. \\ \left. \cup 2D_{(i+b)}^{(2p^j)} \pmod{d_j} \right) \cup \{p^n\}, \text{ and}$$

$$C_1^{(2p^n)} = \bigcup_{j=1}^n \bigcup_{i=0}^{d_j/2-1} p^{n-j} \left( D_{(i+b)}^{(2p^j)} \pmod{d_j} \right. \\ \left. \cup 2D_{(i+b)}^{(2p^j)} \pmod{d_j} \right) \cup \{0\},$$

$$\tilde{C}_0^{(2p^n)} = \bigcup_{j=1}^n p^{n-j} \left( \bigcup_{i=0}^{d_j/2-1} 2D_{(i+b)}^{(2p^j)} \pmod{d_j} \right. \\ \left. \cup \bigcup_{i=d_j/2}^{d_j-1} D_{(i+b)}^{(2p^j)} \pmod{d_j} \right) \cup \{p^n\}, \text{ and}$$

$$\tilde{C}_1^{(2p^n)} = \bigcup_{j=1}^n p^{n-j} \left( \bigcup_{i=0}^{d_j/2-1} D_{(i+b)}^{(2p^j)} \pmod{d_j} \right. \\ \left. \cup \bigcup_{i=d_j/2}^{d_j-1} 2D_{(i+b)}^{(2p^j)} \pmod{d_j} \right) \cup \{0\}. \quad (2)$$

It is obvious that  $\mathbb{Z}_{2p^n} = C_0^{(2p^n)} \cup C_1^{(2p^n)} = \tilde{C}_0^{(2p^n)} \cup \tilde{C}_1^{(2p^n)}$  and  $|C_i^{(2p^n)}| = |\tilde{C}_i^{(2p^n)}| = p^n$ ,  $i = 0, 1$ . Families of balanced binary sequences  $s^\infty = (s_0, s_1, s_2, \dots)$  and  $\tilde{s}^\infty = (\tilde{s}_0, \tilde{s}_1, \tilde{s}_2, \dots)$  of period  $p^n$  can thus be defined as in [6], i.e.,

$$s_i = \begin{cases} 0, & \text{if } i \pmod{p^n} \in C_0^{(2p^n)}, \\ 1, & \text{if } i \pmod{p^n} \in C_1^{(2p^n)}. \end{cases} \quad (3)$$

and

$$\tilde{s}_i = \begin{cases} 0, & \text{if } i \pmod{p^n} \in \tilde{C}_0^{(2p^n)}, \\ 1, & \text{if } i \pmod{p^n} \in \tilde{C}_1^{(2p^n)}. \end{cases} \quad (4)$$

In the case of  $f = 2^r$ , the linear complexity of  $s^\infty, \tilde{s}^\infty$  was estimated in [6], where a conjecture about the linear complexity of these sequences was also made as follows.

**Conjecture.** (1) If  $2^e \equiv -1 \pmod{p}$  but  $2^e \not\equiv -1 \pmod{p^2}$ , then the linear complexity  $L(s^\infty) = 2p^n - (p - 1)$ .

(2) If  $2^e \equiv 1 \pmod{p}$  but  $2^e \not\equiv 1 \pmod{p^2}$ , then the linear complexity  $L(\tilde{s}^\infty) = 2p^n - (p - 1) - e$ .

### 2.3 Main Result

This subsection will study the linear complexity of  $s^\infty, \tilde{s}^\infty$  in (3) and (4) for some even integers  $f$ . The main result in this paper is given as follows.

**Theorem 1** *Let  $p = ef + 1$  be an odd prime with  $2p^{e-1} \not\equiv 1 \pmod{p^2}$  and  $f$  is an even positive integer. Let  $\text{ord}_p(2)$  denote the order of 2 modulo  $p$  and  $v = \text{gcd}(\frac{p-1}{\text{ord}_p(2)}, f)$ .*

*(i) Let  $s^\infty$  be a generalized cyclotomic binary sequence of period  $p^n$  defined in (3). Then the linear complexity of  $s^\infty$  is given by*

$$L(s^\infty) = 2p^n - r \cdot \text{ord}_p(2),$$

where  $0 \leq r \leq \frac{p-1}{\text{ord}_p(2)}$ .

Furthermore, the linear complexity

$$L(s^\infty) = \begin{cases} 2p^n - p + 1, & \text{if } v = f/2; \\ 2p^n, & \text{if } v = 1 \text{ or } 2v \mid \frac{f}{2}, \text{ or } f = v. \end{cases}$$

(ii) Let  $\tilde{s}^\infty$  be a generalized cyclotomic binary sequence of period  $p^n$  defined in (4). Then for the linear complexity of  $\tilde{s}^\infty$  we have

$$2p^n - 2r \cdot \text{ord}_p(2) \leq L(\tilde{s}^\infty) \leq 2p^n - r \cdot \text{ord}_p(2),$$

where  $0 \leq r \leq \frac{p-1}{\text{ord}_p(2)}$ . Furthermore, the linear complexity

$$L(\tilde{s}^\infty) = \begin{cases} 2p^n - 3(p-1)/2 & \text{if } v = f; \\ 2p^n, & \text{if } v \mid \frac{f}{2}, \text{ or } v = 2, v \neq f. \end{cases}$$

**Corollary 2** Let  $f = 2^r$ . Then:

(i) The linear complexity of  $s^\infty$  is given by

$$L(s^\infty) = \begin{cases} 2p^n - p + 1, & \text{if } v = f/2; \\ 2p^n, & \text{otherwise.} \end{cases}$$

(ii) The linear complexity of  $\tilde{s}^\infty$  is given by

$$L(\tilde{s}^\infty) = \begin{cases} 2p^n - 3(p-1)/2, & \text{if } v = f; \\ 2p^n, & \text{otherwise.} \end{cases}$$

**Remark 3** Suppose  $2 \equiv g^u \pmod{p}$  for some integer  $u$ . It is easily seen that  $\text{gcd}(\frac{p-1}{\text{ord}_p(2)}, f) = \text{gcd}(u, f)$ . Thus the condition  $2^e \equiv 1 \pmod{p}$  in Conjecture from [6] is equivalent to  $v = \text{gcd}(\frac{p-1}{\text{ord}_p(2)}, f) = f$  and the condition  $2^e \equiv -1 \pmod{p}$  is equivalent to  $v = f/2$ . In the case that  $f = 2^r$  for a positive integer  $r$ , the integer  $v$  is also a power of 2, which either equals  $f$  or  $f/2$  or divides  $f/4$ . Hence Conjecture from [6] is included in Theorem 1 as a special case. Here we make the correction of Conjecture (ii).

If 2 is a primitive roots modulo  $p$  then  $v = 1$ .

For the proof of Theorem 1 we will use the same definitions and same method that as [4].

Let  $S(x) = s_0 + s_1x + \dots + s_{2p^n-1}x^{2p^n-1}$  and  $\tilde{S}(x) = \tilde{s}_0 + \tilde{s}_1x + \dots + \tilde{s}_{2p^n-1}x^{2p^n-1}$  for the generalized cyclotomic sequences  $s^\infty, \tilde{s}^\infty$  defined in (3) and (4), respectively. Then,

$$S(x) = \sum_{t \in C_1^{(p^n)}} x^t \text{ and } \tilde{S}(x) = \sum_{t \in \tilde{C}_1^{(p^n)}} x^t \quad (5)$$

For simplicity of presentation, we define polynomials as in [4]

$$E_i^{(p^j)}(x) = \sum_{t \in D_i^{(p^j)}} x^t, \quad 1 \leq j \leq n, 0 \leq i < d_j, \quad (6)$$

and

$$H_k^{(p^j)}(x) = \sum_{i=0}^{d_j/2-1} E_{i+k}^{(p^j)} \pmod{d_j}(x), \quad 0 \leq k < d_j, \\ T_k^{(p^m)}(x) = \sum_{j=1}^m H_k^{(p^j)}(x^{p^{m-j}}), \quad m = 1, 2, \dots, n. \quad (7)$$

Notice that the subscripts  $i$  in  $D_i^{(p^j)}, H_i^{(p^j)}(x)$  and  $T_i^{(p^j)}(x)$  are all taken modulo the order  $d_j$ . In the rest of this paper the modulo operation will be omitted when no confusion can arise.

Let  $\overline{\mathbb{F}}_2$  be an algebraic closure of  $\mathbb{F}_2$  and  $\alpha_n \in \overline{\mathbb{F}}_2$  be a primitive  $p^n$ -th root of unity. Denote  $\alpha_j = \alpha_n^{p^{n-j}}, j = 1, 2, \dots, n-1$ .

The properties of considered polynomials were studied in [4]. We have here the following statement.

**Lemma 4** [4] For any  $a \in D_k^{(p^j)}$ , we have

- (i)  $T_i^{(p^m)}(\alpha_m^{p^l a}) = T_{i+k}^{(p^{m-l})}(\alpha_{m-l}) + (p^l - 1)/2 \pmod{2}$  for  $0 \leq l < m$ ; and
- (ii)  $T_i^{(p^m)}(\alpha_m^a) + T_{i+d_m/2}^{(p^m)}(\alpha_m^a) = 1$ .
- (iii) Let  $p$  be a non-Wieferich prime. Then  $T_i^{(p^m)}(\alpha_m) \notin \{0, 1\}$  for  $m > 1$ .
- (iv) Let  $p$  be a non-Wieferich prime. Then  $T_i^{(p^m)}(\alpha_m) + T_{i+f/2}^{(p^m)}(\alpha_m) \neq 1$  for  $m > 1$ .

Throughout this paper an integer  $u$  will be such that  $2 \equiv g^u \pmod{p^n}$ . Now we will show that the studying of linear complexity of above sequences is equivalent to the investigation of properties of  $T_i^{(p^m)}(x)$

**Proposition 5** Let  $\alpha_n$  be a  $p^n$ -th primitive root of unity and let  $2 \equiv g^u \pmod{p^n}$ . Given any element  $a \in \mathbb{Z}_{p^n}$ , we have

- (i)  $S(\alpha_n^a) = 1 + T_b^{(p^n)}(\alpha_n^a) + T_{b+u}^{(p^n)}(\alpha_n^a)$ ; and
- (ii)  $S(\alpha_n^a) = T_b^{(p^n)}(\alpha_n^a) + T_{b+u}^{(p^n)}(\alpha_n^a)$ .

**Proof:** (i) Since  $\sum_{t \in p^{n-j}D_{(i+b)}^{(p^j)}} \alpha_n^{at} = \sum_{t \in p^{n-j}D_{(i+b)}^{(p^j)}} \alpha_n^{at}$  by (1), it follows from our definitions and Lemma 4 that

$$S(\alpha_n^a) = 1 + T_b^{(p^n)}(\alpha_n^a) + T_{b+u}^{(p^n)}(\alpha_n^a).$$

(ii) Similarly we have

$$\tilde{S}(\alpha_n^a) = T_b^{(p^n)}(\alpha_n^a) + T_{b+u}^{(p^n)}(\alpha_n^a). \quad \square$$

We now examine the value of  $T_b^{(p^n)}(\alpha_n^i) + T_{b+u}^{(p^n)}(\alpha_n^i)$  for some integers  $i \in \mathbb{Z}_{p^n}$ .

**Proposition 6** Let  $p$  be a non-Wieferich prime. Then  $S(\alpha_n^i) \neq 0$  and  $\tilde{S}(\alpha_n^a) \neq 0$  for  $i \in \mathbb{Z}_{p^n} \setminus p^{n-1}\mathbb{Z}_p$ .

**Proof:** This is sufficient to prove that  $T_b^{(p^n)}(\alpha_n^i) + T_{b+u}^{(p^n)}(\alpha_n^i) \notin \{0, 1\}$  for  $i \in \mathbb{Z}_{p^n} \setminus p^{n-1}\mathbb{Z}_p$  and  $b = 0, 1, \dots, d_n - 1$ . As it was shown in [4] that without loss of generality it is enough proof,  $T_0^{(p^m)}(\alpha_m) + T_u^{(p^m)}(\alpha_m) \notin \{0, 1\}$  for  $m > 1$ .

We consider two cases.

1. Let  $T_0^{(p^m)}(\alpha_m) + T_u^{(p^m)}(\alpha_m) = 0$ . Since  $(T_0^{(p^m)}(\alpha_m))^2 = T_u^{(p^m)}(\alpha_m) = 0$ , we see that in this case  $T_0^{(p^m)}(\alpha_m) \in \{0, 1\}$ . We obtain a contradiction with Lemma 4 (iii).

2. Let  $T_0^{(p^m)}(\alpha_m) + T_u^{(p^m)}(\alpha_m) = 1$ .

It then follows from Lemma 4 (i) that  $T_{iu}^{(p^m)}(\alpha_m) + T_{(i+1)u}^{(p^m)}(\alpha_m) = 1$  for any integer  $i \geq 1$ . Hence  $T_0^{(p^m)}(\alpha_m) = T_{2iu}^{(p^m)}(\alpha_m)$ .

Denote  $w = \gcd(2u, d_m)$ . Since  $p$  is a non-Wieferich prime, it follows by [4] that  $w$  divides  $f$ . Since the subscript of  $T_i^{(p^m)}(x)$  is taken modulo  $d_m$ , it is easily seen that

$$T_0^{(p^m)}(\alpha_m) = T_{iw}^{(p^m)}(\alpha_m), \quad \text{for any integer } i \geq 1. \tag{8}$$

By Lemma 4 (ii) from the last formula we have  $T_{d_m/2}^{(p^m)}(\alpha_m) = T_{d_m/2+iw}^{(p^m)}(\alpha_m)$  or  $T_{d_m}^{(p^m)}(\alpha_m) = T_{d_m/2+jf}^{(p^m)}(\alpha_m)$ . Then we get that  $T_{d_m/2}^{(p^m)}(\alpha_m) = T_{f/2}^{(p^m)}(\alpha_m)$ . Thus, by Lemma 4 (ii) we obtain that  $T_0^{(p^m)}(\alpha_m) + 1 = T_{f/2}^{(p^m)}(\alpha_m)$ . But the latest equality is not possible for  $m > 1$  by Lemma 4 (iv).  $\square$

By Proposition 6, we only need to study the value of  $T_b^{(p^n)}(\alpha_n^i) + T_{b+u}^{(p^n)}(\alpha_n^i)$  for integers  $i$  in the set  $p^{n-1}\mathbb{Z}_p$ . Suppose  $i = p^{n-1}a$ ,  $a \in D_i^{(p)}$ . Then, it follows from Proposition 5 and Lemma 4 that

$$S(\alpha_n^i) = 1 + H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1),$$

where  $k \equiv b + i \pmod{f}$ . The following proposition examines the value of  $H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1)$  according to the relation between  $f$  and  $\text{ord}_p(2)$ .

**Proposition 7** Let  $p = ef + 1$  be an odd prime with  $f$  being an even positive integer and  $v = \gcd(\frac{p-1}{\text{ord}_p(2)}, f)$ . Then,

$$(i) \quad \left| \left\{ k \in \mathbb{Z}_f \mid H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1) = 0 \right\} \right| = \begin{cases} f, & \text{if } v = f, \\ 0, & \text{if } v|f/2 \text{ or } v = 2, v \neq f. \end{cases}$$

$$(ii) \quad \left| \left\{ k \in \mathbb{Z}_f \mid H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1) = 1 \right\} \right| = \begin{cases} f, & \text{if } v = f/2, \\ 0, & \text{if } v = 1, \text{ or } v = f \text{ or } 2v|f/2. \end{cases}$$

**Proof:** Since  $\text{ord}_p(2) = \frac{p-1}{\gcd(p-1, u)}$ , it follows that  $\gcd(u, f) = \gcd(\frac{p-1}{\text{ord}_p(2)}, f) = v$  [4].

(i) For  $v = f$  this statement is clear.

Let  $v|f/2$  or  $v = 2, v \neq f$ . We shall prove this case by contradiction. Suppose  $H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1) = 0$  for some integer  $k$ . Since  $(H_k^{(p)}(\alpha_1))^2 = H_{k+u}^{(p)}(\alpha_1)$ , it follows that  $H_k^{(p)}(\alpha_1) \in \{0, 1\}$ . By [4] this is not possible for  $v|f/2$  or  $v = 2, v \neq f$ .

(ii) For  $v = f/2$  this statement is clear. If  $v = f$  then  $2 \in D_0^{(p)}$  and we have  $H_k^{(p)}(\alpha_1) + H_k^{(p)}(\alpha_1) = 1$ . This is impossible

Suppose  $H_k^{(p)}(\alpha_1) + H_{k+u}^{(p)}(\alpha_1) = 1$  for some integer  $k$ . Without loss of generality, we assume  $k = 0$  and  $H_0^{(p)}(\alpha_1) = H_u^{(p)}(\alpha_1) + 1$ .

In the case when  $v \neq f$ . Since  $\gcd(u, f) = \gcd(\frac{p-1}{\text{ord}_p(2)}, f) = v$ , by a similar argument as in the proof of Proposition 6 we get

$$H_0^{(p)}(\alpha_1) = H_{2v}^{(p)}(\alpha_1) = \dots = H_{2vi}^{(p)}(\alpha_1).$$

So, if  $2v$  divides  $f/2$ , then  $H_{f/2}^{(p)}(\alpha_1) = H_{2v \cdot f/4v}^{(p)}(\alpha_1) = H_0^{(p)}(\alpha_1)$ , which is a contradiction.

Let  $v = 1$ . Then we get  $H_i^{(p)}(\alpha_1) + H_{i+1}^{(p)}(\alpha_1) + 1 = 0, i = 0, 1, \dots, f - 1$  and then  $E_i^{(p)}(\alpha_1) + E_{i+f/2}^{(p)}(\alpha_1) + 1 = 0, i = 0, 1, \dots, f - 1$ . In [4] it was shown that this is impossible.  $\square$

**Proof of Theorem 1.** Recall that the linear complexity of  $s^\infty$  is given by

$$L(s^\infty) = N - \deg \left( \gcd \left( (x^{p^n} - 1)^2, S(x) \right) \right).$$

(i) From Proposition 6 we know  $S(\alpha_n^i) \neq 0$  for  $i \in \mathbb{Z}_{p^n} \setminus p^{n-1}\mathbb{Z}_p$ . For the remaining set  $p^{n-1}\mathbb{Z}_p$ , if  $i = 0$ , then  $S(1) = 1$ ; if  $i \in p^{n-1}\mathbb{Z}_p^*$ , we have

$$S(\alpha_n^i) = 1 + H_b^{(p)}(\alpha_1^a) + H_{b+u}^{(p)}(\alpha_1^a)$$

for some integer  $a \in \mathbb{Z}_p^*$ .

Suppose  $H_k^{(p)}(\alpha_1^a) + H_{k+u}^{(p)}(\alpha_1^a) = 1$  for some integer  $k$ . Then

$$1 = (H_k^{(p)}(\alpha_1))^2 + H_{k+u}^{(p)}(\alpha_1^2) = H_{k+u}^{(p)}(\alpha_1) + H_{k+2u}^{(p)}(\alpha_1),$$

and so on (here  $u \not\equiv 0 \pmod{f}$ ). So, we have

$$|\{i : S(\alpha_n^i) = 0, i = 1, 2, \dots, p^n - 1\}| = r \text{ord}_p(2).$$

where  $r$  is an integer with  $0 \leq r \leq \frac{p-1}{\text{ord}_p(2)}$ .

Further, by (5) we see that

$$xS'(x) = \sum_{j=1}^n \sum_{i=0}^{d_j/2-1} \sum_{t \in D_{i+b}^{(2p^j)} \pmod{d_j}} x^{p^{n-j}t}.$$

Hence,  $\alpha_n^i S(\alpha_n^i) = T_b^{(p^n)}(\alpha_n^i)$ . So, if  $\alpha_n^i$  is a root of  $S(x)$  and  $S'(x)$  then  $1 + T_b^{(p^n)}(\alpha_n^i) + (T_b^{(p^n)}(\alpha_n^i))^2 = 0$  and  $T_b^{(p^n)}(\alpha_n^i) = 0$ . It is not possible and any root of  $S(x)$  is simple.

Then the statement of this theorem follows from Proposition 6.

(ii) In this case

$$S(\alpha_n^i) = H_b^{(p)}(\alpha_1^a) + H_{b+u}^{(p)}(\alpha_1^a)$$

for some integer  $a \in \mathbb{Z}_p^*$ .

Then as earlier we again get

$$|\{i : S(\alpha_n^i) = 0, i = 1, 2, \dots, p^n - 1\}| = r \text{ord}_p(2).$$

where  $r$  is an integer such that  $0 \leq r \leq \frac{p-1}{\text{ord}_p(2)}$ .

Here, by (5) we see that

$$x\tilde{S}'(x) = \sum_{j=1}^n \sum_{i=0}^{d_j/2-1} \sum_{t \in D_{i+b}^{(2p^j)} \pmod{d_j}} x^{p^{n-j}t}.$$

and also  $\alpha_n^i \tilde{S}(\alpha_n^i) = T_b^{(p^n)}(\alpha_n^i)$ . If  $v = f$  then it follows from [4] that

$$|\{i : T_b^{(p^n)}(\alpha_n^i) = 0, i = 1, 2, \dots, p^n - 1\}| = (p-1)/2.$$

Then the statement of this theorem follows from Proposition 6. □

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