Criterions of supersinguliarity and groups of Montgomery and Edwards curves in cryptography

RUSLAN SKURATOVSKII Interregional Academy of Personnel Management, Kyiv, and NTU of Igor Sikorsky, Kiev, UKRAINE

VOLODYMYR OSADCHYY Dept. of computer science UCF, Florida, USA

Abstract—We consider the algebraic affine and projective curves of Edwards over the finite field Fpn. It is well known that many modern cryptosystems can be naturally transformed into elliptic curves. The criterions of the supersingularity of Montgomery and Edwards curves are found. In this paper, we extend our previous research into those Edwards algebraic curves over a finite field and we construct birational isomorphism of them with cubic in Weierstrass normal form. One class of twisted Edwards is researched too. We propose a novel effective method of point counting for both Edwards and elliptic curves. In addition to finding a specific set of coefficients with corresponding field characteristics for which these curves are supersingular, we also find a general formula by which one can determine whether or not a curve Ed[Fp] is supersingular over this field. The method proposed has complexity O($p \log^2_2 p$). This is an improvement over both Schoof's basic algorithm and the variant which makes use of fast arithmetic (suitable for only the Elkis or Atkin primes numbers) with complexities O($\log^{4}_{2} pn$) and O($\log^{4}_{2} pn$) respectively. The embedding degree of the supersingular curve of Edwards over Fpn in a finite field is additionally investigated. Singular points of twisted Edwards curve are completely described. Due existing the birational isomorphism between twisted Edwards curve and elliptic curve in Weierstrass normal form. Also it is considered minimum degree of an isogeny (distance) between curves of this two classes when such isogeny exists. We extend the existing isogenous of elliptic curves.

Keywords: finite field, elliptic curve, Edwards curve, group of points of an elliptic curve, criterion of supersingularity of curves, order of curve counting method.order of curve counting method.

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1. Introduction

One of the main approaches of post-quantum cryptography (PQC) is using of the isogeny of supersingular elliptic curves as well as Edwards curves which studied by us in [10], [11] with as many subgroups of their points as possible. The security of the SIDH algorithm [2] requires that the number of subgroups of the curve E_d of the order $p+1 = 4 \cdot 3^n 5^m$ for protection against a quantum computer has to be more than 760 bits. To reach this goal we investigate conditions of supersingularity of Montgomery and Edwards curves. Moreover we present our criterions of supersingularity of Montgomery and Edwards curve.

In modern cryptography, supersingular curves find more and more applications, in particular in identity-based cryptosystems and for short group signatures [31], [33]–[35], [37], in PQC. At the same time, they have numerous advantages over their competitors that are MNT curves [36].

In order to construct a cryptosystem based on an elliptic curve, we need to firstly analyze the order of a group of elliptic curve (EC) points. We provides an approach to construct Edwards curves of determined order, which are important within the cryptography and coding theory domains. It should be noted that it was accepted in 1999 as an ANSI standard and in 2000 as an IEEE and NIST standard. One of the fundamental problems in EC cryptography regards the generation of cryptographically secure ECs over prime fields, suitable for use in various cryptographic applications. Because supersingular elliptic curves are vulnerable to pairing-based attacks, we find a criterion for Edwards curve supersingularity [10]. The method of finding the order of an algebraic curve over a finite field F_{p^n} is of huge interest currently and is at the center of many mathematical studies connected with the use of groups of points of curves of genus 1. In this article, this critical problem is solved. All proofs and analytical results belong to Skuratovskii R. V. but computational examples, confirming statements, are made together with other authors.

Our algorithm has improved complexity for algebraic extensions with a large degree over finite fields. This is because upon choosing sufficiently large values n, we obtain $\mathcal{O}(\log_2^8 p^n)$, which has a much larger complexity than $\mathcal{O}(p \log_2^2 p)$ for some fixed p.

We denote the order of projective Edwards curve with coefficient d over finite field F_{p^n} by $\overline{N}_{d[p^n]}$ and order of affine Edwards curve by $N_{d[p^n]}$.

2. Materials and Methods

Definition. The isogeny of the elliptic curve E(K) over a field K into a curve E'(K) is a homomorphism p : $E(K) \to E'(K)$ over an algebraic closure K given by rational functions. This means that for all points $P, Q \in$ $E(K), \phi(P+Q) = \phi(P) + \phi(Q)$ and there exist rational functions

$$\phi(x,y) = (\frac{p(x)}{q(x)}, y\frac{f(x)}{g(x)}) = (x', y')$$

mapping points of the curve E at the points of the curve E'.

The degree of isogeny is the maximum of the degrees $a = deg\phi(x, y) = max\{degp(x), degq(x)\}$, and its kernel is subgroup $N \subseteq E$ of the order α (separable isogeny), the points of which are mapped by the function $\phi(x, y)$ into a neutral element of the group O.

Isogeny compresses the points of the curve E at α times and is a surjection (α points of the curve E are mapped to one point of the curve E'). When N = 0, isogeny becomes isomorphism (a = 1).

The method of finding the order of an algebraic curve over a finite field \mathbb{F}_{p^n} is clearly related with constructing curves of a given order. We propose a method for counting points from Edwards and elliptic curves in direct response to a paper by Schoof [8]. To calculate the sum of the squares of binomial coefficients in a finite modulus, we make use of the recursive calculation of the factorial by multiplication. The method we propose has improved complexity over both Schoof's basic algorithm and the variant which makes use of fast arithmetic (suitable only for Elkis or Atkin primes numbers) with complexities $\mathcal{O}(\log_2^8 p^n)$ and $\mathcal{O}(\log_2^4 p^n)$ respectively.

In particular, our method yields a vastly improved complexity for algebraic extensions with large degree n of finite fields. This occurs because when choosing sufficiently large values n, Schoof's algroithm has complexity $O(\log_2^8 p^n)$ [8], which clearly is much larger than $O(p \log_2^2 p)$ for fixed p.

We now recall an important result from Vinogradov [17] which will act as criterion for supersingularity.

Lemma 1: Let $k \in \mathbb{N}$ and $p \in \mathbb{P}$. Then

$$\sum_{k=1}^{p-1} k^n \equiv \begin{cases} 0 \pmod{p}, & n \not| (p-1), \\ -1 \pmod{p}, & n | (p-1), \end{cases}$$

where n|(p-1) means that n is divisible by p-1.

We additionally propose a method for counting the points from Edwards curves and elliptic curves in form of Montgomery in response to an earlier paper by Schoof [8].

Let *d* denote some large number. The method of Karatsuba multiplication is used to calculate all the values $d^j \pmod{p}$. In this task, it is optimally applying recursive multiplication d^{j-1} upon *d* and for this we utilise the Karatsuba multiplication method, which requires $\mathcal{O}(\log_2^{\log_2 3} p)$, rather than apply the Barrett method of modular multiplication. The complexity of computing the entire tuple of degrees d^j , $j = 1, \ldots, n$ is therefore $\mathcal{O}(\frac{p-1}{2}\log_2^{\log_2 3} p)$.

3. Review of Previous Results

The method of finding supersingular Edwards curves was partially studied by us in [10], [11], [32], [35]–[37]. One of

the main goals of post-quantum cryptography (PQC) is the isogeny of supersingular elliptic curves which studied by us in [10], [11] with as many subgroups of their points as possible.

It is known that the theoretical results of S. Stepanov [16] and P. Deligne provide both upper and lower bounds for the number of curve points, that is $p^n + 1 - \omega_1^n - \omega_2^n$ [25]. The complexity of the fastest algorithms known are Schoof's basic algorithm [8], which yields $\mathcal{O}(\log_2^8 p^n)$, as well as a variant that makes use of fast arithmetic (suitable only for Elkis or Atkin primes), which has complexity $\mathcal{O}(\log_2^4 p^n)$. Our method is faster than the approach to order curve determination by counting of a longest chain of points using dividing of a point on 2 [23].

4. Systematic Algebraic Analysis qh'tj g'Ewt xg'cpf 'Ewt xg'Order Calculation Method

The *twisted Edwards curve* with coefficients $a, d \in F_p^*$, is the curve $E_{a,d}$:

$$ax^{2} + y^{2} = 1 + dx^{2}y^{2}, (1)$$

where ad(a - d) = 0, d = 1, p = 2. It should be noted that a twisted Edwards curve is simply called an *Edwards curve* when a = 1. By E_d , we denote the Edwards curve with coefficient $d \in F_p^*$ defined as

$$x^2 + y^2 = 1 + dx^2 y^2,$$

over \mathbb{F}_p . The projective curve has the form

$$F(x, y, z) = ax^{2}z^{2} + y^{2}z^{2} = z^{4} + dx^{2}y^{2}.$$

The special points are the infinitely distant points (1,0,0)and (0,1,0) and we find its singularities at infinity in the corresponding affine components $A^1 := az^2 + y^2z^2 = z^4 + dy^2$ and $A^2 := ax^2z^2 + z^2 = z^4 + dx^2$. These are entitled the simple singularities.

The condition of singular points [4], [24] give us the following system:

$$\begin{cases} \frac{F(x,y)}{\partial x} = 0, \\ \frac{F(x,y)}{\partial y} = 0, \end{cases} \begin{cases} 2ax = 2dxy^2, \\ 2y = 2dx^2y, \end{cases} \begin{cases} ax - dxy^2 = 0, \\ y - dx^2y = 0, \end{cases} \\ \begin{cases} x(a - dy^2) = 0, \\ y(1 - dx^2) = 0, \end{cases} \\ \begin{cases} x = 0, \\ (a - dy^2) = 0, \\ 1 - dx^2 = 0, \\ y = 0, \end{cases} \begin{bmatrix} (0,0) \notin E_{a,d}, \\ (\pm \sqrt{\frac{1}{d}}, 0) \\ (0, \pm \sqrt{\frac{a}{d}}) \end{cases}$$

The choice of pairs of values in coordinates is justified by the fact that the point $(0,0) \notin E_{a,d}$ The images of points $\left(\pm\sqrt{\frac{1}{d}},0\right)$. and $\left(0,\pm\sqrt{\frac{a}{d}}\right)$ belongs to projective curve as singular points with coordinates $\left(\pm\sqrt{\frac{1}{d}},\infty\right)$ and $\left(\infty,\pm\sqrt{\frac{a}{d}}\right)$. We describe the structure of the local ring at the point

we describe the structure of the local ring at the point p_1 , whose elements are quotients of functions with the form $F(x, y, z) = \frac{f(x, y, z)}{g(x, y, z)}$, where the denominator cannot take value 0 at the singular point p_1 . In particular, it should be

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noted that a local ring with two singularities consists of those functions with denominators which are not divisible by (x-1)(y-1).

We denote by $\delta_p = \dim \frac{\mathcal{O}_p}{\mathcal{O}_p}$, where \mathcal{O}_p denotes the local ring at the singular point p generated by the relations of regular functions

$$\mathcal{O}_p = \left\{ \frac{f}{g} : (g, (x-1)(y-1)) = 1 \right\}$$

and $\overline{\mathcal{O}}_p$ denotes the whole closure of the local ring at the singular point p.

We find that $\delta_p = 1$ is the dimension of the factor as a vector space because the basis of extension $\frac{\overline{O}_p}{\overline{O}_p}$ consists of just one element at each distinct point. We then calculate the genus of the curve according to Fulton [4], namely

$$\rho^*(C) = \rho_{\alpha}(C) - \sum_{p \in E} \delta_p = \frac{(n-1)(n-2)}{2} - \sum_{p \in E} \delta_p$$

$$\rho^*(C) = 3 - 2 = 1,$$

where $\rho_{\alpha}(C)$ denotes the arithmetic genus of the curve C with parameter $n = \deg(C) = 4$. It should be noted for completeness that the supersingular points were discovered in [10]. Recall that the curve has a genus of 1 and in consequence, it is known to be isomorphic to a flat cubic curve. Despite this, the curve is importantly not elliptic because of its singularity in the projective part.

Both the Edwards curve and the twisted Edwards curve are isomorphic to some affine part of the elliptic curve.

Koblitz [4], [5] tells us that one can detect if a curve is supersingular using the search for the curve when that curve has the same number of points as its torsion curve. In addition, an elliptic curve E over F_q is called supersingular if for every finite extension F_{q^r} , there are no points in the group $E(F_{q^r})$ of order p [20]. It is known [1] that the transition from an Edwards curve to the related torsion curve is determined by the reflection $(\overline{x}, \overline{y}) \mapsto (x, y) = (\overline{x}, \frac{1}{\overline{y}})$.

The *order of a curve* is precisely the number of its affine points with a neutral element, where the group operation is well defined.

Here under x^0 we understand the identity 1.

Let a sum of the form $\sum_{x=0}^{p-1} P(x)$, where P(x) is a polynomial of degree at most 2p - 2. In case of the Montgomery curve we will have maximal degree $\frac{3p-3}{2}$. Let be $P(x) = a_{2p-2}x^{a_{2p-2}} + a_{2p-3}x^{a_{2p-3}} + \ldots + a_1x + a_0 = \sum_{k=0}^{2p-2} a_k x^k$. Then $\sum_{x=0}^{p-1} P(x) = \sum_{x=0}^{p-1} \sum_{k=0}^{2p-2} a_k x^k = \sum_{k=0}^{2p-2} (a_k \cdot \sum_{x=0}^{p-1} x^k)$. Let be $\sum_{x=0}^{p-1} x^k = S_k$, according to Lemma 1,

$$S_{k} \equiv \begin{cases} 0(modp), \ if \ k = 0 \ or \ k \ \not(p-1) \\ -1(modp), \ if \ k \ \not 0 \ and \ k \ (p-1) \end{cases}$$

It means $\sum_{x=0}^{p-1} P(x) \equiv \sum_{k=0}^{2p-2} a_k S_k \equiv -a_{p-1} - a_{2p-2}(modp)$

We recall the basic Montgomery form for an elliptic curve

$$By^2 = x^3 + Ax^2 + x (2)$$

Statement : The Montgomery curve (2) is supersingular over F_p if and only if

$$\sum_{j=0}^{\frac{p-1}{2}} (C^j_{\frac{p-1}{2}})^2 r^{\frac{p-1}{2}-2j} \equiv 0 (modp),$$

where r is one of the roots of the equation $x^2 + Ax + 1 = 0$. Note that degree of $r^{\frac{p-1}{2}-2j}$ can be reduced to r^{-2j} due to Fermat's theorems, by multiplication the previous congruence on power $r^{\frac{p-1}{2}}$ because of p is prime.

To prove this statement we need to find the number N, which is equal to the number of points on the curve above F_p

$$By^{2} = x(x - r)(x - c).$$
 (3)

Also we need to research the conditions when $N \equiv p + 1$. Let us take into account that from the most square equation obtained on the right-hand side it follows that the roots are mutually inverse, i.e. $c = r^{-1}$. Suppose that the roots r, cquadratic trinomial $x^2 + Ax + 1 = 0$ belong to the field F_p . Since the roots divide 1, they are reciprocal, i.e. kind r, r^{-1} therefore it is more convenient to work with the polynomial

$$By^{2} = x(x-r)(x-r^{-1})$$
(4)

but not (3) $By^2 = x(x-r)(x-c)$ the left side is a square residue once $\left(\frac{B}{p}\right) = 1$. By Euler's theorem, we can determine the quadraticity of the right part of (3) as

$$By^{2} = x^{\frac{p-1}{2}} (x-r)^{\frac{p-1}{2}} (x-r^{-1})^{\frac{p-1}{2}}$$
(5)

For a fixed value x the number of solutions of equation (3) is equal to

 $1 + \left(\frac{x(x-r)(x-r^{-1})}{p}\right), \text{ where } \left(\frac{a}{p}\right) \text{ is symbol of Legendre.}$ In case $\left(\frac{B}{p}\right) = -1$, then the last amount takes the form $1 - \left(\frac{x(x-r)(x-r^{-1})}{p}\right)$ but further magnitudes $\left(\frac{x(x-r)(x-r^{-1})}{p}\right)$ absolutely the same. According to Euler's formula $\left(\frac{a}{p}\right) = a^{\frac{p-1}{2}} (modp)$, therefore for a fixed x the number of solutions of equation (3) is compared modulo p from $1 + (x(x - r)(x - c^{-1}))^{\frac{p-1}{2}}$. The maximum coefficient in (5) is $\frac{3p-3}{2}$ therefore the coefficient a_{2p-2} at 2p - 2 degree is equal to 0. Then we calculate the sums $\left(\frac{x(x-r)(x-c)}{p}\right)$ and the corresponding amounts $1 + \left(\frac{x(x-r)(x-c)}{p}\right)$ also quite similar and $1 - \left(\frac{x(x-r)(x-c)}{p}\right)$ to all $x \in F_p$. So, summing up all x, we have $N = \sum_{x=0}^{p-1} 1 + (x(x - r)(x - r^{-1}))^{\frac{p-1}{2}} \equiv \sum_{x=0}^{p-1} x^{\frac{p-1}{2}} (x-r)^{\frac{p-1}{2}} (x-r^{-1})^{\frac{p-1}{2}} modp$. To calculate equal need N coefficient at degree $\frac{(p-1)}{2}$. But further we will study the internal sum no longer by x and in powers of Newton's $(\sum_{j=0}^{p-1} C_{\frac{p-1}{2}}^j r^{j} x^{j} (r^{-1})^{\frac{p-1}{2}-j}) (modp)$. Now we will turn to indexing both sums by j in order to successfully multiply them

$$N \equiv x^{\frac{p-1}{2}} \left(\sum_{j=0}^{p-1} C_{\frac{p-1}{2}}^{j} x^{j} (-r)^{\frac{p-1}{2}-j} \right)$$
$$\left(\sum_{j=0}^{p-1} C_{\frac{p-1}{2}}^{j} x^{\frac{p-1}{2}-j} (-r^{-1})^{j} \right) (modp).$$

From these parentheses we choose the power of the variable x, which is equal p-1 and adding them we get the coefficient at x^{p-1} . General form of the coefficient at the degree p-1 is such

$$\begin{split} &(-1)^{\frac{p-1}{2}}C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}(r)^{\frac{p-1}{2}-j}C_{\frac{p-1}{2}}^{j}(r^{-1})^{j}x^{p-1} = \\ &(-1)^{\frac{p-1}{2}}C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}r^{\frac{p-1}{2}-j}C_{\frac{p-1}{2}}^{j}r^{-j}x^{p-1} = \\ &= (-1)^{\frac{p-1}{2}}C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}C_{\frac{p-1}{2}}^{j}r^{\frac{p-1}{2}-2j}x^{p-1}. \end{split}$$

Taking in consideration, for our method, the inequality p > 2, then the whole amount of coefficients near degree of x^{p-1} in (5) is such

$$\begin{split} &\sum_{j=0}^{p-1} C_{\frac{p-1}{2}}^{p-1} -j} (-r)^{\frac{p-1}{2}-j} C_{\frac{p-1}{2}}^{j} r^{-j} x^{p-1} \quad \text{notice, that} \\ &C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j} = C_{\frac{p-1}{2}}^{j}. \text{ Therefore, the previous amount is} \\ &\text{simplified } (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^{j})^2 r^{\frac{p-1}{2}-2j} x^{p-1} \text{ and we get} \\ &\text{that } a_{p-1} \equiv (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^{j})^2 r^{\frac{p-1}{2}-2j} (modp). \\ &\text{Thus, the equality} \end{split}$$

$$a_{p-1} \equiv \sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 r^{\frac{p-1}{2}-2j} \equiv 0 (modp)$$

provides N = p + 1, which had to be proved.

To compute the first *n* lines of Pascal's triangle by (modp), you need $O(n^2 log_2(p))$ Squaring by the Karatsuba method spend $O(\log_2^{\log_2 3} p)$ operations. But we will compute Binomial coefficient more quickly by recursive way.

We will now strengthen an existing result given in [10], [12]–[16], [24]. Let $N_{d[p]}$ denote the number of points with a neutral element of an affine Edwards curve over the finite field \mathbb{F}_p and let $\overline{N}_{d[p]}$ denote the number of points on the projective curve over the same field.

Theorem 1: If
$$p \equiv 3 \pmod{4}$$
 is prime and

$$\sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j \equiv 0 \pmod{p},$$
(6)

is true, then the orders of the curves $x^2 + y^2 = 1 + dx^2y^2$ and $x^2 + y^2 = 1 + d^{-1}x^2y^2$ over F_p are equal to

$$N_{d[p]} = p + 1$$

if
$$\left(\frac{d}{p}\right) = -1$$
, and $N_{d[p]} = p - 3,$

if $\left(\frac{d}{p}\right) = 1$.

Proof: Consider the Edwards curve E_d , namely

$$x^2 + y^2 = 1 + dx^2 y^2. (7)$$

We transform (7) into the form $y^2(1 - dx^2y^2) = 1 - x^2$ and then we express y^2 by applying a rational transformation, which leads to $y^2 = \frac{1-x^2}{1-dx^2y^2}$. For technical reasons, we transform this into the curve

$$y^{2} = (x^{2} - 1)(dx^{2} - 1).$$
(8)

Let $M_{d[p]}$ denote the number of points from an affine Edwards curve over the finite field \mathbb{F}_p . The curve (8) has precisely $M_{d[p]} = N_{d[p]} + \left(\frac{d}{p}\right) + 1$ points. This is exactly $\left(\frac{d}{p}\right) + 1$ greater than the number of points of E_d , where $\left(\frac{d}{p}\right)$ denote the Legendre Symbol. Let $a_0, a_1, \ldots, a_{2p-2}$ be the coefficients of the polynomial $a_0 + a_1x + \cdots + a_{2p-2}x^{2p-2}$, which was obtained from $(x^2 - 1)^{\frac{p-1}{2}}(dx^2 - 1)^{\frac{p-1}{2}}$ after opening the brackets.

Upon summing over all x, we yield that

$$M_{d[p]} = \sum_{x=0}^{p-1} 1 + ((x^2 - 1)(dx^2 - 1))^{\frac{p-1}{2}} =$$
$$= p + \sum_{x=0}^{p-1} (x^2 - 1)^{\frac{p-1}{2}} (dx^2 - 1)^{\frac{p-1}{2}} \equiv$$
$$\equiv \sum_{x=0}^{p-1} (x^2 - 1)^{\frac{p-1}{2}} (dx^2 - 1)^{\frac{p-1}{2}} (\mod p).$$

By opening the brackets in $(x^2-1)^{\frac{p-1}{2}}(dx^2-1)^{\frac{p-1}{2}}$, we have that $a_{2p-2} = (-1)^{\frac{p-1}{2}} \cdot d^{\frac{p-1}{2}} \equiv \left(\frac{d}{p}\right) \pmod{p}$. In light of Lemma 1, we have

$$M_{d[p]} \equiv -\left(\frac{d}{p}\right) - a_{p-1} \pmod{p}.$$
(9)

Recall that we need to prove that $M_{d[p]} \equiv 1 \pmod{p}$ if $p \equiv 3 \pmod{8}$ and $M_{d[p]} \equiv -1 \pmod{p}$ if $p \equiv 7 \pmod{8}$. Therefore, we have to show that $M_{d[p]} \equiv -(\frac{d}{p}) - a_{p-1} \pmod{8}$. p for $p \equiv 3 \pmod{4}$ if $\sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j \equiv 0 \pmod{p}$. It should be noted that if we prove that $a_{p-1} \equiv 0 \pmod{p}$, then the result will follow from (8). We now determine a_{p-1} using Newton's binomial formula. In particular, the coefficient of x^{p-1} in the polynomial, namely a_{p-1} , is obtained from the product $(x^2-1)^{\frac{p-1}{2}} (dx^2-1)^{\frac{p-1}{2}}$. Therefore, $a_{p-1} = (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} d^j (C_{\frac{p-1}{2}}^j)^2$. Instead of this equality, we notice that the following equality holds:

$$\sum_{j=0}^{\frac{p-1}{2}} d^j (C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}) (-1)^{\frac{p-1}{2}-(\frac{p-1}{2}-j)} \cdot d^j (C_{\frac{p-1}{2}}^j)^2 (-1)^{\frac{p-1}{2}-j}$$

$$= (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} d^j C^{\frac{p-1}{2}-j}_{\frac{p-1}{2}} \cdot C^j_{\frac{p-1}{2}} = (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} d^j (C^j_{\frac{p-1}{2}})^2.$$

Noting that $a_{p-1} = -\sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j$ means that the exact number of affine points on a non-supersingular curve is

$$M_{d[p]} \equiv -a_{2p-2} - a_{p-1} \equiv -\left(\frac{d}{p}\right) + \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^{2} d^{j} \pmod{p}.$$
(10)

In accordance with the condition of this theorem, we have $a_{p-1} = 0$ and therefore $M_{d[p]} \equiv -a_{2p-2} \pmod{p}$. In consequence, if p is a prime such that $p \equiv 3 \pmod{4}$ and $\sum_{j=0}^{p-1} (C_{\frac{p-1}{2}}^j)^2 d^j \equiv 0 \pmod{p}$ holds, then the curve E_d has

$$N_{d[p]} = p - \left(\frac{d}{p}\right) - \left(\left(\frac{d}{p}\right) + 1\right) = p - 1 - 2\left(\frac{d}{p}\right)$$

affine points and a group of points of the curve completed by singular points has p + 1 points.

The exact number of the points yields an upper bound of 2p + 1 since each $x \neq 0$ corresponds two values of y, while, for x = 0, we have only that y = 0. Taking into account that $x \in F_p$, we have exactly p values of x. In addition, there are 4 pairs $(\pm 1, 0)$ and $(0, \pm 1)$, which are points of E_d , thus $N_{d[p]} > 1$. In consequence, $N_{d[p]} = p + 1$, which completes the proof.

Corollary 1: The orders of the curves $x^2 + y^2 = 1 + dx^2y^2$ and $x^2 + y^2 = 1 + d^{-1}x^2y^2$ over F_p are equal to

$$N_{d[p]} = p + 1 = \overline{N}_{d[p]},$$

if $\left(\frac{d}{n}\right) = -1$, and

$$N_{d[p]} = p - 3 = \overline{N}_{d[p]} - 4,$$

if $(\frac{d}{p}) = 1$, if and only if p is a prime with $p \equiv 3 \pmod{4}$ and

$$\sum_{j=0}^{\frac{p-1}{2}} (C^j_{\frac{p-1}{2}})^2 d^j \equiv 0 \pmod{p}.$$

Theorem 2: If the coefficient d = 2 or $d = 2^{-1}$ and $p \equiv 3 \pmod{4}$, then

$$\sum_{j=0}^{\frac{p-1}{2}} d^j (C^j_{\frac{p-1}{d}})^2 \equiv 0 \pmod{p} \text{ and } \overline{N}_{d[p]} = p+1.$$

 After applying the congruence $\left(\frac{p-1}{2}-k\right)^2 \equiv \left(\frac{p-1}{2}+1+k\right)^2 \pmod{p}$ with $0 \le k \le \frac{p-1}{2}$ to the multipliers in previous parentheses, we obtain $\left[\left(\frac{p-1}{2}\right)\left(\frac{p-1}{2}-1\right)\cdots(j+1)\right]$. It yields

$$\left(\frac{p-1}{2}\right) \left(\frac{p-1}{2}-1\right) \cdots \left(\frac{p-1}{2}-j+1\right) \left[\left(\frac{p-1}{2}+1\right) \cdots \left(\frac{p-1}{2}+\frac{p-1}{2}-j\right) \right] (-1)^{\frac{p-1}{2}-j}.$$

Thus, as a result of squaring, we have:

$$\left(\left(\frac{p-1}{2}\right)!C_{\frac{p-1}{2}}^{j}\right)^{2} \equiv \left(\frac{p-1}{2}-j+1\right)^{2}.$$

$$\left(\frac{p-1}{2}-j+2\right)^{2}\cdots\left(p-j-1\right)^{2}(\text{mod }p)$$
(11)

It remains to prove that

$$\sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^j\right)^2 2^j \equiv 0 \pmod{p},$$

if $p \equiv 3 \pmod{4}$.

Consider the auxillary polynomial

$$P(t) = \left(\frac{p-1}{2}!\right)^2 \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^j\right)^2 t^j.$$

We are going to show that P(2) = 0 and therefore $a_{p-1} \equiv 0 \pmod{p}$. Using (11) it can be shown that $a_{p-1} = P(t) = \left(\frac{p-1}{2}!\right)^2 \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^j\right)^2 t^j \equiv \sum_{j=0}^{\frac{p-1}{2}} (k+1)^2 (k+2)^2 \cdots \left(\frac{p-1}{2}+k\right)^2 t^k \pmod{p}$ over F_p .

We replace d by t in (6) such that we can research a more generalised problem. It should be noted that $P(t) = \partial^{\frac{p-1}{2}} \left(\partial^{\frac{p-1}{2}} (Q(t) t^{\frac{p-1}{2}}) t^{\frac{p-1}{2}} \right)$ over F_p , where $Q(t) = t^{p-1} + \cdots + t + 1$ and $\partial^{\frac{p-1}{2}}$ denotes the $\frac{p-1}{2}$ -th derivative by t, where t is new variable but not a coordinate of curve. Observe that $Q(t) = \frac{t^p-1}{t-1} \equiv \frac{(t-1)^p}{t-1} \equiv (t-1)^{p-1} \pmod{p}$ and therefore the following equality

$$P(t) = \left(\left((t-1)^{p-1} t^{\frac{p-1}{2}} \right)^{\left(\frac{p-1}{2}\right)} t^{\frac{p-1}{2}} \right)^{\left(\frac{p-1}{2}\right)} \text{ holds over } F_p.$$

In order to simplify notation we let $\theta = t-1$ and $R(\theta) =$

In order to simplify notation we let $\theta = t - 1$ and $R(\theta) = P(\theta + 1)$. For the case t = 2 we have $\theta = 1$. Performing this substitution leads the polynomial P(t) of 2 to the polynomial R(t - 1) of 1. Taking into account the linear nature of the substitution $\theta = t - 1$, it can be seen that that derivation by θ and t coincide. Derivation leads us to the transformation of polynomial $R(\theta)$ to form where it has the necessary coefficient a_{p-1} . Then

$$\begin{aligned} R(\theta) &= P(\theta+1) = \\ &= \partial^{\frac{p-1}{2}} \left(\partial^{\frac{p-1}{2}} \left(\theta^{p-1} (\theta+1)^{\frac{p-1}{2}} \right) (\theta+1)^{\frac{p-1}{2}} \right) = \\ &= \partial^{\frac{p-1}{2}} \left(\frac{(p-1)!}{((p-1)/2)!} \theta^{\frac{p-1}{2}} (\theta+1)^{\frac{p-1}{2}} \right). \end{aligned}$$

In order to prove that $a_{p-1} \equiv 0 \pmod{p}$, it is now sufficient to s how that $R(\theta) = 0$ if $\theta = 1$ over F_p . We obtain

$$R(1) = \frac{(p-1)!}{(\frac{p-1}{2})!} \sum_{j=0}^{\frac{p-1}{2}} C^{j}_{\frac{p-1}{2}}(j+1) \cdots (j+\frac{p-1}{2}).$$
(12)

We now will manipulate with the expression $(\frac{p-1}{2} - j + 1)(\frac{p-1}{2} - j + 2)\cdots(\frac{p-1}{2} - j + \frac{p-1}{2})$. In order to illustrate the simplification we now consider the scenario when p = 11 and hence $\frac{p-1}{2} = 5$. The expression gets the form $(5-j+1)(5-j+2)\cdots(5-j+5) = (6-j)(7-j)\cdots(10-j) \equiv ((-5-j)(-4-j)\cdots(-1-j)) \equiv (-1)^5((j+1)(j+2)\cdots(j+5)) \pmod{11}$. Therefore, for a prime p, we can rewrite the expression as $(\frac{p-1}{2} - j + 1)(\frac{p-1}{2} - j + 2)\cdots(\frac{p-1}{2} - j + \frac{p-1}{2}) \equiv (-1)^{\frac{p-1}{2}}(j+1)\cdots(j+\frac{p-1}{2}) \equiv -1(j+1)\cdots(j+\frac{p-1}{2}) \pmod{p}$.

As a result, the symmetrical terms in (12) can be reduced yielding $a_{p-1} \equiv 0 \pmod{p}$. It should be noted that $(-1)^{\frac{p-1}{2}} = -1$ since p = Mk + 3 and $\frac{p-1}{2} = 2k + 1$. Consequently, we have P(2) = R(1) = 0 and hence $a_{p-1} \equiv 0 \pmod{p}$ as required. Thus, $\sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 \equiv 0 \pmod{p}$, completing the proof of the theorem.

Corollary 2: The curve E_d is supersingular iff $E_{d^{-1}}$ is supersingular.

Proof: Let us recall the proved fact in Theorem IV that
$$N_{d[p]} \equiv -a_{2p-2} - a_{p-1} \equiv -\left(\frac{d}{p}\right) + \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j \pmod{p}.$$
Since $\left(C_{\frac{j}{2}}^{j}\right)^2 d^j \equiv 0 \pmod{p}$ by condition and the congru-

Since $(C_{\frac{p-1}{2}}^{j})^2 d^j \equiv 0 \pmod{p}$ by condition, and the congruence $(\frac{d}{p}) \equiv (\frac{d^{-1}}{p})$ holds, then $N_{d[p]} \equiv N_{d^{-1}[p]}$. Now we estimate the number of points on the curve (8). Let

Now we estimate the number of points on the curve (8). Let $M_{d[p]}$ denote the number of solutions to equation (8) over the field F_p . It should be observed that for x = 1 and x = -1, the right side of (8) is equal to 0. Due to this the number $M_{d[p]}$ can therefore be bounded by

$$2 \le M_{d[p]} \le 2p - 2,$$
 (13)

where if $a_{p-1} \equiv 0 \pmod{p}$ we have $N_{d[p]} \equiv -\left(\frac{d}{p}\right) \pmod{p}$. Note that the number of solutions is bounded by $N_{d[p]} \leq 2p-2$ because if x = 1 and x = -1 we only have one value of y, namely y = 0. For different values of x, we will have no more than two solutions for y because the equation (8) is quadratic relative to y. Thus, the only possible number is $M_{d[p]} \equiv p - \left(\frac{d}{p}\right) \pmod{p}$.

$$\begin{split} M_{d[p]} &\equiv p - \left(\frac{d}{p}\right) \pmod{p}.\\ Corollary 3: \text{ If } p \equiv 3 \pmod{4}, \text{ is prime then there exists}\\ \text{some } T \text{ such that } T &\equiv \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j \leq 2\sqrt{q} \text{ and } N_{d[p]} = \\ p - 1 - 2\left(\frac{d}{p}\right) + T. \end{split}$$

Proof: Due to equality (10) and the bounds (13) as well as according to generalized Hasse-Weil theorem $|N_{d[p]} - (p + 1) - 2\left(\frac{d}{p}\right)| \le 2g\sqrt{p}$, where g is genus of curve, we obtain exact number $N_{d[p]}$. As we showed, g = 1. From Theorem IV as well as from Corollary 2 we get, that $\sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j \equiv$ $-N_{d[p]} - (p+1) - 2\left(\frac{d}{p}\right)$ so there exists $T \in \mathbb{Z}$, such that $T < 2\sqrt{p}$ and $N_{d[p]} = p - 1 - 2\left(\frac{d}{p}\right) + T$. Example1: If p = 13, d = 2 gives $N_{2[13]} = 8$ and p = 13, $d^{-1} = 7$ gives that the number of points of E_7 is $N_{7[13]} = 20$. Moreover, if $p \equiv 7 \pmod{8}$, then the order of torsion subgroup of curve is $N_2 = N_{2^{-1}} = p - 3$, which is clearly different to p + 1.

For instance p = 31, then $N_{2[31]} = N_{2^{-1}[31]} = 28 = 31 - 3$, which is clearly not equal to p + 1. If $p = 7, d = 2^{-1} \equiv (4 \mod 7)$ then the curve $E_{2^{-1}}$ has four points, namely (0, 1); (0, 6); (1, 0); (6, 0), and the in case p = 7 with $d = 2 \pmod{7}$, the curve $E_{2^{-1}}$ also has four points: (0, 1); (0, 6); (1, 0); (6, 0), demonstrating the order in this scenario is p - 3.

The following theorem shows that the total number of affine points upon the Edwards curves E_d and $E_{d^{-1}}$ are equal under certain assumptions. This theorem additionally provides us with a formula for enumerating the number of affine points upon the birationally isomorphic Montgomery curve N_M .

Theorem 3: Let d satisfy the condition of supersingularity (6). If $n \equiv 1 \pmod{2}$ and p is prime, then

$$\overline{N}_{d[p^n]} = p^n + 1,$$

and the order of curve is equal to

$$N_{d[p^n]} = p^n - 1 - 2\left(\frac{d}{p}\right).$$

If $n \equiv 0 \pmod{2}$ and p is prime, then the order of curve

 $N_{d[p^n]} = p^n - 3 - 2(-p)^{\frac{n}{2}},$

and the order of projective curve is equal to

$$\overline{N}_{d[p^n]} = p^n + 1 - 2(-p)^{\frac{n}{2}}.$$

Proof: We consider the extension of the base field F_p to F_{p^n} in order to determine the number of the points on the curve $x^2 + y^2 = 1 + dx^2y^2$. Let P(x) denotes a polynomial with degree m > 2 whose coefficients are from \mathbb{F}_p . To make the proof, we take into account that it is known [16] that the number of solutions to $y^2 = P(x)$ over \mathbb{F}_{p^n} will have the form $p^n + 1 - \omega_1^n - \cdots - \omega_{m-1}^n$ where $\omega_1, \ldots, \omega_{m-1} \in \mathbb{C}$, $|\omega_i| = p^{\frac{1}{2}}$.

In case of our supersingular curve, if $n \equiv 1 \pmod{2}$ the number of points on projective curve over \mathbb{F}_{p^n} is determined by the expression $p^n + 1 - \omega_1^n - \omega_2^n$, where $\omega_i^n \in \mathbb{C}$ and $\omega_1 = -\omega_2$, $|\omega_i| = \sqrt{p}$ that's why $\omega_1 = i\sqrt{p}$, $\omega_2 = -i\sqrt{p}$ with $i \in \{1, 2\}$. In the general case, it is known [6], [16] that $|\omega_i| = p^{\frac{1}{2}}$. The order of the projective curve is therefore $p^n + 1$.

If $p \equiv 7 \pmod{8}$, then it is known from a result of Skuratovskii [10] that E_d has in its projective closure of the curve singular points which are not affine and therefore

$$N_{d[p]} = p^n - 3$$

If $p \equiv 3 \pmod{8}$, then there are no singular points, hence

$$\overline{N}_{d[p]} = N_{d[p]} = p^n + 1.$$

Consequently the number of points on the Edwards curve depends on $\left(\frac{d}{p}\right)$ and is equal to $N_{d[p]} = p^n - 3$ if $p \equiv 7 \pmod{8}$ and $N_{d[p]} = p^n + 1$ if $p \equiv 3 \pmod{8}$ where $n \equiv 1 \pmod{2}$. We note that this is because the transformation of (8) in E_d depends upon the denominator $(dx^2 - 1)$.

If $n \equiv 1 \pmod{2}$ then, with respect to the sum of root of of the characteristic equation for the Frobenius endomorphism $\omega_1^n + \omega_2^n$, which in this case have the same signs, we obtain that the number of points in the group of the curve is $p^n + 1 - \omega_1^n - \omega_2^n$ [25].

For $n \equiv 0 \pmod{2}$ we always have, that every $d \in F_p$ is a quadratic residue in F_{p^n} . Consequently, because of $\left(\frac{d}{p}\right) = 1$ four singular points appear on the curve. Thus, the number of affine points is less by 4, i.e. $N_{d[p^n]} = p^n - 1 - 2\left(\frac{d}{p}\right) - 2(-p)^{\frac{n}{2}} = p^n - 3 - 2(-p)^{\frac{n}{2}}$. In more details ω_1 , ω_2 are eigen values of the Frobenius operator F endomorphism on etale cohomology over the finite field \mathbb{F}_{p^n} , where F acts of $H^i(X)$. The number of points, in general case, are determined by Lefshitz formula:

$$#X(\mathbf{F}_{p^n}) = \sum (-1)^i tr(F^n \mid H^i(X)),$$

where $\#X(\mathbb{F}_{p^n})$ is a number of points in the manifold X over \mathbb{F}_{p^n} , F^n is composition of Frobenius operator. In our case, E_d is considered as the manifold X over \mathbb{F}_{p^n} .

We continue to study birational isomorphism $x = \frac{1+u}{1-u}$, $y = \frac{2u}{v}$ which we proposed in [11] between E_d and E_M . In the following Lemma it will be extended from E_M on $E_{a,d}$.

Lemma 2: There exists birational isomorphism of E_d also $E_{a,d}$ with E_M and with elliptic curve in canonical Weierstrass form $E_{a,b}$, which is determined by correspondent mappings $x = \frac{1+u}{1-u}, y = \frac{2u}{v}$ and Velu formulas [29], over finite field F_{p^n} .

Proof: To verify this statement in both the supersingular case and the non-supersingular case and in also for both values of the quadratic residue $\frac{d}{p}$ we the birational equivalence $(u,v) \mapsto (2u/v, (u-1)/(u+1)) = (x,y)$ between E_M and E_d . In supersingular case if $n \equiv 1 \pmod{2}$ according to Theorem 3 we know that the curve

$$x^2 + y^2 = 1 + dx^2 y^2$$

over F_p contains $p-1-2\left(\frac{d}{p}\right)$ points (x, y), with coordinates over prime field F_p .

In supersingular case according to Theorem 3 if $n \equiv 1 \pmod{2}$ and p is prime then the curve E_d over F_{p^n} has order

$$N_{d[p^n]} = p^n - 1 - 2\left(\frac{d}{p}\right).$$

If $n \equiv 0 \pmod{2}$ and p is prime, then the order of curve

$$N_{d[p^n]} = p^n - 3 - 2(-p)^{\frac{n}{2}}$$

Consider the transformation of the curve $x^2 + y^2 = 1 + dx^2y^2$ into the following form $y^2(dx^2 - 1) = x^2 - 1$. Make the substitutions $x = \frac{1+u}{1-u}$ and $y = \frac{2u}{v}$. We will call the

special points of this transformations the point in which these transformations or inverse transformations are not determined. As a result the equation of curve become in the form

$$\frac{4u^2}{v^2} \cdot \frac{(d-1)u^2 + 2(d+1)u + (d-1)}{(1-u)^2} = \frac{4u}{(1-u)^2}.$$

Multiply the equation of the curve by

$$\frac{v^2(1-u)^2}{4u}$$

As a result of the reduction, we obtain the equation

$$v^{2} = (d-1)u^{3} + 2(d+1)u^{2} + (d-1)u.$$

We analyze what new solutions appeared in the resulting equation in comparing with $y^2(dx^2 - 1) = x^2 - 1$.

First, there is an additional solution (u, v) = (0, 0). Second, if d is a quadratic residue by modulo p, then the solutions appear

$$(u_1, v_1) = \left(\frac{-(d+1) - 2\sqrt{d}}{d-1}, 0\right),$$
$$(u_2, v_2) = \left(\frac{-(d+1) + 2\sqrt{d}}{d-1}, 0\right).$$

If $\left(\frac{d}{p}\right) = -1$ then as it was shown above the order of E_d is equal to p + 1. Therefore, in case $\left(\frac{d}{p}\right) = -1$ order of E_d appears one additional solution of from (u, 0) more exact it is point with coordinates (0, 0) also two points ((-1; 0), (1; 0)) of E_d have not images on E_M in result of action of birational map on E_M . Thus, in this case, number of affine points on E_M is equal to p + 1 - 2 + 1 = p. The table of correspondence between points is the following.

-
-
-
-
-
(-1, 0)
(1, 0)

TABLE I SPECIAL POINTS OF BIRATIONAL MAPING

If x = -1 then equality $x = \frac{1+u}{1-u}$ transforms to form -1 + u = 1 + u, or -1 = 1 that is impossible for p > 2. therefore point (-1, 0) have not an image on E_M .

Consider the case x = 1. As a result of the substitutions x = (1+u)/(1-u), y = 2u/v we get the pair (x, y) corresponding to the pair (u, v) for which $v^2 = (d-1)u^3 + 2(d+1)u^2 + (d-1)u$.

If it occurs that y = 0, then the preimage having coordinates u = 0 and v is not equal to 0 is suitable for the birational map $y = \frac{2u}{v}$ which implies that u = 0 and $v \neq 0$. But pair (u, v) of such form do not satisfies the equation of obtained in result

of mapping equation of Montgomery curve $v^2 = (d-1)u^3 +$ $2(d+1)u^2 + (d-1)u$. Therefore the corresponding point (u, v) will not be a solution to the equation $v^2 = (d-1)u^3 +$ $2(d+1)u^2 + (d-1)u$, since there will be an element on the left side, different from 0, and on the right will be 0. That is a contradiction as required, therefore (x, y) = (1, 0) is the special point having not image on E_M .

If y = 0 then in equality $y = \frac{2u}{v}$ appear zeros in numerator

and denominator and transformation is not correct. The points $(\frac{-(d+1)-2\sqrt{d}}{d-1}, 0)$, $(\frac{-(d+1)+2\sqrt{d}}{d-1}, 0)$, $(1, -2\sqrt{d})$, $(1, 2\sqrt{d})$ exist on E_M only when $(\frac{d}{p}) = 1$. These points are elements of group which can be presented on Riemann sphere over F_q . The points $(1, -2\sqrt{d})$, $(1, 2\sqrt{d})$ have not images on E_d because of in denominator of transformations $x = \frac{1+u}{1-u}$ appears zero. By the same reason points $\left(\frac{-(d+1)-2\sqrt{d}}{d-1},0\right)$, $(\frac{-(d+1)+2\sqrt{d}}{d-1},0)$ have not an images on $E_d.$

If $\left(\frac{d}{p}\right) = 1$ then as it was shown above the order of E_d is equal to p-3. Therefore order of E_M is equal to p because of 5 additional solutions of equation of E_M appears but 2 points ((-1;0),(1;0)) of E_d have not images on E_M . These are 5 additional points appointed in tableau above. Also it exists one infinitely distant point on an Montgomery curve. Thus, the order of E_M is equal p+1 in this case as supersingular curve has.

It should be noted that the supersingular curve E_d is birationally equivalent to the supersingular elliptic curve which may be presented in Montgomery form $v^2 = (d-1)u^3 + 2(d+1)u^3 +$ $1)u^2 + (d-1)u$. As well as exceptional points [1] for the birational equivalence $(u, v) \mapsto (2u/v, (u-1)/(u+1)) = (x, y)$ are in one to one correspondence to the affine point of order 2 on E_d and to the points in projective closure of E_d . Since the formula for number of affine points on E_M can be applied to counting $N_{d[p]}$. In such way we apply this result [7], [16], to the case $y^2 = P(x)$, where degP(x) = m, m = 3. The order $N_{M[p^n]}$ of the curve E_M over F_{p^k} can be evaluated due to Stepanov [16]. The research tells us that the order is $N_{M[p^n]} = p^n + 1 - \omega_1^n - \omega_2^n$, where $\omega_i^n \in \mathbb{C}$ and $\omega_1^n = -\omega_2^n$, $|\omega_i| = \sqrt{p}$ with $i \in \{1, 2\}$. Therefore, we conclude when $n \equiv 1 \pmod{2}$, we know the order of Montgomery curve is precisely $N_{M[p^n]} = p^n + 1$. This result leads us to the conclusion that the number of solutions of $x^2 + y^2 = 1 + dx^2y^2$ as well as $v^2 = (d-1)u^3 + 2(d+1)u^2 + (d-1)u$ over the finite field \mathbb{F}_{p^n} are determined by the expression $p^n + 1 - \omega_1^n - \omega_2^n$ if $n \equiv 1 \pmod{2}$.

Consider a more general case, namely the transition from a twisted Edwards curve to E_M .

Firstly the twisted Edwards curve $E_{a,d}$ by a rational transformation

$$\phi(x,y) = \left((a-d)\frac{1+u}{1-v}, (a-d)\frac{2u}{y} \right) = ((a-d)x, (a-d)y) = (u,v).$$
(14)

is transformed into the birationally equivalent Montgomery form

$$v^{2} = u^{3} - 2(a+d)u^{2} + (a-d)u.$$
(15)

The point (0.0) is the second-order point of this curve, which, together with the point at infinity as a neutral element of the group, forms the kernel of the 2-isogeny. It is required to find parameters \bar{a} and d of the isogenous curve with equation (15) and the rational function 2(x, y) = (X, Y). For the Montgomery curve in the general form

$$E_{M(c,b)}: y^2 = x^3 + cx^2 + bx, \tag{16}$$

finding 2-isogeny is well known [29].

To construct the isomorphic mapping from E_m to the elliptic curve in normal form of Weierstrass for p > 3, $y^2 = x^3 + ax + b$ we apply linear substitution to coordinates of Montgomery curve points: $u = x - \frac{c}{3}$, v = y. This linear substitution yields the following transformation:

$$\begin{split} v^2 &= \left(x - \frac{c}{3}\right)^3 + c\left(x - \frac{c}{3}\right)^2 + \left(x - \frac{c}{3}\right) = \\ &= x^3 - x^2C + x\frac{C^2}{3} - \frac{C^3}{27} + x^2C - 2x\frac{C^2}{3} + \frac{C^3}{9} + x - \frac{C}{3} = \\ &= x^3 + x\frac{C^2}{3} - 2x\frac{C^2}{3} + x - \frac{C}{3} + \frac{C^3}{9} - \frac{C^3}{27} = \\ &= x^3 + \left(\frac{C^2}{3} - 2\frac{C^2}{3} + 1\right)x + \left(\frac{2C^3}{27} - \frac{C}{3}\right) = \\ &= x^3 + \left(1 - \frac{C^2}{3}\right)x + \left(\frac{2C^3}{27} - \frac{C}{3}\right) = x^3 + ax + b. \end{split}$$

Thus, we obtain the presentation of elliptic curve coefficients $a = 1 - \frac{C^2}{3}$, $b = \frac{2C^3}{27} - \frac{C}{3}$ For the Montgomery curve of the general form (16)

$$\begin{split} v^2 &= \left(x - \frac{c}{3}\right)^3 + c\left(x - \frac{c}{3}\right)^2 + e\left(x - \frac{c}{3}\right) = \\ &= x^3 - x^2C + x\frac{C^2}{3} - \frac{C^3}{27} + x^2C - 2x\frac{C^2}{3} + \frac{C^3}{9} + ex - \frac{eC}{3} = \\ &= x^3 + x\frac{C^2}{3} - 2x\frac{C^2}{3} + x - \frac{C}{3} + \frac{C^3}{9} - \frac{C^3}{27} = \\ &= x^3 + \left(\frac{C^2}{3} - 2\frac{C^2}{3} + e\right)x + \left(\frac{2C^3}{27} - \frac{eC}{3}\right) = \\ &= x^3 + \left(e - \frac{C^2}{3}\right)x + \left(\frac{2C^3}{27} - \frac{eC}{3}\right) = x^3 + ax + b. \end{split}$$

Based on the Velu formulas [27], [29] which were also mentioned and used in [28], using the laws of the addition of the points of the curve in the general Weierstrass form, for the curve (16) one can obtain the 2-isogeny ([29]. the example 12.4)

$$\psi(u,v) = \left(\frac{u^2 + cu + b}{u}, \frac{u^2 - b}{u^2}v\right) = (X,Y)$$
(17)

as a result the equation of the isogenous curve is the following:

$$Y^{2} = X^{3} - 2cX^{2} + (c^{2} - 4b)X.$$
 (18)

The discriminant of the quadratic equation $X^2 - 2cX + c$ $(c^2 - 4b) = 0$ obtained from the right-hand side of (18) is $\delta = 16b$, and depending on the meaning of $(\frac{b}{p})$, the curve (18) has one or three points of the 2^{nd} order. In the first case, one can construct one 2-isogeny. in the two-three points (for three kernels as subgroups of the second-order).

We propose to extend the isogenous (17) up to cubic in normal form of Weierstrass without second power by the new isogenous,

$$\psi(u,v) = \left(\frac{u^2 + cu + b}{u} + \frac{2c}{3}, \frac{u^2 - b}{u^2}v\right) = (X,Y)$$
(19)

then new obtain an equation $Y^2 = X^3 + dX + c$. In fact we put $X + \frac{2c}{3}$ instead of X and as we showed above where we applied substitutions $u = x - \frac{c}{3}$, v = y in E_M , this linear substitution reduces second power in equation of E_M but free constant b arises.

For twisted Edwards curve $E_{a,d}$ as follows from equations of (15) and (16), only those curves $E_{a,d}$ of general (canonical) form can be reduced to the Montgomery form (E_M) (and, accordingly, to the Edwards form), the parameter b of which is the quadratis square $(\frac{b}{p}) = 1$). This is connected with the existence on the curve canonical elliptic curve in the form (15) the points with the order 4 i.e. $F = (u_1, v_1)$, such that 4F = (0, 0). Then, taking $b = u_1^2$, equation of $E_{M,(c,b)}$ (16) after the substitution $c \to Cu_1$ is reduced to the form

$$v^2 = u^3 + Cu_1 u^2 + u_1^2 u. (20)$$

or to isomorphic (or its quadratic torsion) curve (1)

$$v^{2} = u^{3} + Cu^{2} + u, C = 2\frac{a+d}{a-d}$$
(21)

This curve is birationally equivalent to the generalized Edwards form curve (1) when $v^2 \rightarrow Dv^2$. The equation (20) is equivalent to (15) when $u_1^2 = (a-d)^2$ and $Cu_1 = 2(a+d)$.

Thus, the 2-isogenic curve (18) with the discriminant $\delta = 16u_1^2$ in this case, has three points of the 2^{nd} order, and corresponding isogenies can be found only in the classes of quadratic Edwards curve with $\left(\frac{a}{p}\right) = 1$, $\left(\frac{d}{p}\right) = 1$ which is completely described in [28], and twisted Edwards curves forming pairs of quadratic twist. At the time the curve $E_{a,d}$, for which isogeny is constructed, can have one point of the 2^{nd} order and two points of the 4^{th} order (the class of Edwards curves $E_{a,d}$ with $\left(\frac{ad}{p}\right) = -1$ which is completely described in [28]), or belong to other classes of Edwards curves with three points of the second order.

The inverse transformation of isogenous curves after the substitution $c \to Cx_1$ in $E_{a,d}$ in the Montgomery form into the Edwards form [28], [30] $E_{a,ad}$ is performed based on rational functions this Lemma taking into account different values of coordinates of points of the 4^{th} order $\pm X_1 \in 4a\sqrt{d}, 4a\sqrt{1-d}, 4a\sqrt{d(d-1)}$ or $\pm X_1 = \bar{a} - \bar{d}$ with the help of rational function.

Example: For instance, in case $p \equiv 3mod4$ supersingular elliptic curve $v^2 = u^3 + u$ (for which $\left(\frac{(c^2-4b)}{p}\right) = -1$) has one point of the second-order and two points of the 4 order and is isomorphic to the Edwards curve E_d , where $\left(\frac{d}{p} = -1\right)$.

Example: The elliptic curve presented in the form of Montgomery E_M : $v^2 = u^3 + 6u^2 + u$, is birationally equivalent [1] to the curve $x^2 + y^2 = 1 + 2x^2y^2$ over the field F_{p^k} .

Corollary 4: If d = 2, $n \equiv 1 \pmod{2}$ and $p \equiv 3 \pmod{8}$, then the order of curve E_d and order of the projective curve are the following:

$$\overline{N}_{d[p^n]} = p^n + 1.$$

If d = 2, $n \equiv 1 \pmod{2}$ and $p \equiv 7 \pmod{8}$, then the number of points of projective curve is

$$\overline{N}_{d[p^n]} = p^n + 1,$$

and the number of points on affine curve E_d is also

$$N_{d[p^n]} = p^n - 3$$

In case d = 2, $n \equiv 0 \pmod{2}$, $p \equiv 3 \pmod{4}$, the general formula of the curves order is

$$N_{d[p^n]} = p^n - 3 - 2(-p)^{\frac{n}{2}}.$$

The general formula for $n \equiv 0 \pmod{2}$ and d = 2 for the number of points on projective curve for the supersingular case is

$$\overline{N}_{d[p^n]} = p^n + 1 - 2(-p)^{\frac{n}{2}}.$$

Proof: We denote by $N_{M[p^n]}$ the order of the curve E_M over F_{p^n} . The order $N_{M[p^n]}$ of E_M over F_{p^n} can be evaluated [6] as $N_{M[p^n]} = p^n + 1 - \omega_1^n - \omega_2^n$, where $\omega_i^n \in \mathbb{C}$ and $\omega_1^n = -\omega_2^n$, $|\omega_i| = \sqrt{p}$ with $i \in \{1, 2\}$. For the finite algebraic extension of degree n, we will consider $p^n - \omega_1^n - \omega_2^n = p^n$ if $n \equiv 1 \pmod{2}$. Therefore, for $n \equiv 1 \pmod{2}$, the order of the Montgomery curve is precisely given by $N_{M[p^n]} = p^n + 1$. Here's one infinitely remote point as a neutral element of the group of points of the curve.

Considering now an elliptic curve, we have $\omega_1 = \bar{\omega}_2$ by [5], which leads to $\omega_1 + \omega_2 = 0$. For n = 1, it is clear that $N_M = p$. When n is odd, we have $\omega_1^n + \omega_2^n = 0$ and therefore $N_{M,n} = p^n + 1$. Because n is even by initial assumption, we shall show that $N_{M[p^n]} = p^n + 1 - 2(-p)^{\frac{n}{2}}$ holds as required.

Note that for n = 2 we can express the number as $\overline{N}_{d[p^2]} = p^2 + 1 + 2p = (p+1)^2$ with respect to Lagrange theorem have to be divisible on $\overline{N}_{d[p]}$. Because a group of $E_d(F_{p^2})$ over square extension of F_p contains the group $E_d(F_p)$ as a proper subgroup. In fact, according to Theorem 1 the order $E_d(F_p)$ is p+1 therefore divisibility of order $E_d(F_{p^2})$ holds because in our case p = 7 thus $\overline{N}_{E_d} = 8^2$ and $p+1 = 8 = N_{d[7]}$ [19].

The following two examples exemplify Corollary 4. Example: If $p \equiv 3 \pmod{8}$ and n = 2k then we have

when d = 2, n = 2, p = 3 that the number of affine points equals to

$$N_{2[3]} = p^n - 3 - 2(-p)^{\frac{n}{2}} = 3^2 - 3 - 2 \cdot (-3) = 12,$$

and the number of projective points is equal to

$$\overline{N}_{2[3]} = p^n + 1 - 2(-p)^{\frac{n}{2}} = 3^2 + 1 - 2 \cdot (-3) = 16.$$

Example: If $p \equiv 7 \pmod{8}$ and n = 2k then we have when d = 2, n = 2, p = 7 that the number of affine points equals to

$$N_{2[7]} = p^n - 3 - 2(-p)^{\frac{n}{2}} = 7^2 - 3 - 2 \cdot (-7) = 60,$$

and the number of projective points is equal to

$$\overline{N}_{2[7]} = p^n + 1 - 2(-p)^{\frac{n}{2}} = 7^2 + 1 - 2 \cdot (-7) = 64.$$

Proposition 1: The group of points of the supersingular curve E_d contains $p - 1 - 2\left(\frac{d}{p}\right)$ affine points and the affine singular points whose number is $2\left(\frac{d}{p}\right) + 2$.

Proof: The singular points were discovered in [10] and hence if the curve is free of singular points then the group order is p+1.

Example: The number of curve points over finite field when d = 2 and p = 31 is equal to $N_{2[31]} = N_{2^{-1}[31]} =$ p - 3 = 28.

Theorem 4: The order of projective Edwards curve over F_p is congruent to

$$\overline{N}_{d[p]} \equiv \left(p - 1 - 2\left(\frac{d}{p}\right) + (-1)^{\frac{p+1}{2}} \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j\right) \equiv$$

$$\equiv \left((-1)^{\frac{p+1}{2}} \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j} \right)^2 d^j - 1 - 2\left(\frac{d}{p}\right) \right) (\mod p).$$

The true value of $\overline{N}_{d[p]}$ lies in [4;2p] and is even.

Proof: This result follows from the number of solutions of the equation

 $y^2 = (dx^2 - 1)(x^2 - 1)$ over F_p which equals to

$$\sum_{x=0}^{p-1} \left(\frac{(x^2-1)(dx^2-1)}{p}) + 1 \right) \equiv$$

$$\equiv \sum_{x=0}^{p-1} \left(\frac{(x^2 - 1)(dx^2 - 1)}{p} \right) + p \equiv$$

$$\equiv (\sum_{j=0}^{\frac{p-1}{2}} (x^2 - 1)^{\frac{p-1}{2}} (dx^2 - 1)^{\frac{p-1}{2}}) \pmod{p} \equiv$$

$$\equiv \left((-1)^{\frac{p+1}{2}} \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^{2} d^{j} - \left(\frac{d}{p}\right)\right) (\bmod \ p) \,.$$

The quantity of solutions for $x^2 + y^2 = 1 + dx^2y^2$ differs from the quantity of $y^2 = (dx^2 - 1)(x^2 - 1)$ by $(\frac{d}{p}) + 1$ due to

new solutions in the from $(\sqrt{d}, 0)$, $(-\sqrt{d}, 0)$. So this quantity is such

$$\begin{split} &\sum_{x=0}^{p-1} \left(\frac{(x^2-1)(dx^2-1)}{p}) + 1 \right) - \left((\frac{d}{p}) + 1 \right) \equiv \\ &\sum_{x=0}^{p-1} \left(\frac{(x^2-1)(dx^2-1)}{p}) \right) + p - \left((\frac{d}{p}) - 1 \right) \equiv \\ &\equiv \left(\sum_{j=0}^{\frac{p-1}{2}} (x^2-1)^{\frac{p-1}{2}} (dx^2-1)^{\frac{p-1}{2}} - (\frac{d}{p}) + 1 \right) (\bmod \ p) \equiv \\ &\equiv (-1)^{\frac{p+1}{2}} \sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j - (2(\frac{d}{p}) + 1) (\bmod \ p) \,. \end{split}$$

According to Lemma 1, the last sum $\left(\sum_{k=0}^{\frac{p-1}{2}} (x^2-1)^{\frac{p-1}{2}} (dx^2-1)^{\frac{p-1}{2}}\right) \pmod{p}$ is congruent to $-a_{p-1} - a_{2p-2} \pmod{p}$, where a_i are the coefficients from $(x^{2}-1)^{\frac{p-1}{2}}(dx^{2}-1)^{\frac{p-1}{2}} = a_{0} + a_{1}x + \dots + a_{2p-2}x^{2p-2}.$ Last presentation was obtained due to transformation $(x^2 - 1)^{\frac{p-1}{2}} (dx^2 - 1)^{\frac{p-1}{2}} = (\sum_{x=0}^{p-1} C_{\frac{p-1}{2}}^k x^{2k} (-1)^{\frac{p-1}{2}-k}) (\sum_{x=0}^{p-1} C_{\frac{p-1}{2}}^j d^j x^{2j} (-1)^{\frac{p-1}{2}-j}).$ Therefore a_{2p-2} is equal to $d^{\frac{p-1}{2}} \equiv (\frac{d}{p}) \pmod{p}$ and

 $a_{p-1} = \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j} \right)^2 d^j (-1)^{\frac{p-1}{2}}.$

According to Newton's binomial formula a_{p-1} equal to the coefficient at x^{p-1} in the product of two brackets and when substituting it d instead of 2 is such

$$(-1)^{\frac{p-1}{2}}\sum_{j=0}^{\frac{p-1}{2}}d^{j}(C^{j}_{\frac{p-1}{2}})^{2},$$

that is, it has the form of the polynomial with inverse order of coefficients.

Indeed, we have equality

$$\sum_{j=0}^{\frac{p-1}{2}} d^{j} \left(C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}\right) (-1)^{\frac{p-1}{2}-(\frac{p-1}{2}-j)} \cdot \left(C_{\frac{p-1}{2}}^{j}\right)^{2} (-1)^{\frac{p-1}{2}-j} = \\ = (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} d^{j} C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j} \cdot C_{\frac{p-1}{2}}^{j} = (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} d^{j} \left(C_{\frac{p-1}{2}}^{j}\right)^{2}.$$

In form of a sum it is the following

$$\begin{split} &\sum_{j=0}^{\frac{p-1}{2}} 2^j (C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j}) (-1)^{\frac{p-1}{2}-(\frac{p-1}{2}-j)} \cdot 2^j (C_{\frac{p-1}{2}}^j)^2 (-1)^{\frac{p-1}{2}-j} = \\ &= (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} 2^j C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j} \cdot C_{\frac{p-1}{2}}^j = (-1)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} 2^j (C_{\frac{p-1}{2}}^j)^2. \end{split}$$

Now, if

$$a_{p-1} = \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^j\right)^2 d^j (-1)^{\frac{p-1}{2}} \equiv 0 \pmod{p},$$

then as it is was shown by the author in that the curve is supersingular and the number of solutions of the

$$y^2 = (dx^2 - 1)(x^2 - 1)$$

over F_p equals to

$$p - 1 - 2\left(\frac{d}{p}\right) + \left(1 + \left(\frac{d}{p}\right)\right) = p - \left(\frac{d}{p}\right)$$

and differs from the quantity of solutions of

$$x^2 + y^2 = 1 + dx^2 y^2$$

by $\left(\frac{d}{p}\right) + 1$ due to new solutions of

$$y^2 = (dx^2 - 1)(x^2 - 1).$$

Thus, in general case if

$$a_{p-1} = \sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j (-1)^{\frac{p-1}{2}} \neq 0,$$

we have

 \equiv

$$N_{E_d} = \left(p - \left(\frac{d}{p}\right) - \left(\left(\frac{d}{p}\right) + 1\right) - \left(-1\right)^{\frac{p-1}{2}} \cdot \frac{p-1}{2} \cdot \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{\frac{p-1}{2}-j} C_{\frac{p-1}{2}}^{j}\right)^2 d^j\right) \equiv$$
$$\equiv \left(p - 1 - \left(-1\right)^{\frac{p-1}{2}} \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j - 2\left(\frac{d}{p}\right)\right) \equiv$$
$$\left(\left(-1\right)^{\frac{p+1}{2}} \sum_{j=0}^{\frac{p-1}{2}} \left(C_{\frac{p-1}{2}}^{j}\right)^2 d^j - 1 - 2\left(\frac{d}{p}\right)\right) \pmod{p}.$$

The exact order is not less than 4 because cofactor of this curve is 4. To determine the order is uniquely enough to take into account that p and 2p have different parity. Taking into account that the order is even we chose a term p or 2p, for the sum which define the order.

Let us analyze the complexity of calculating the value of

$$\sum_{j=0}^{\frac{p-1}{2}} (C_{\frac{p-1}{2}}^j)^2 d^j.$$

Binomial coefficients of the form $C_{\frac{p-1}{2}}^{l}$ we calculate recursively having $C_{\frac{p-1}{2}}^{l}$ we get $C_{\frac{p-1}{2}}^{l+1}$. Such a transformation can be done by one multiplication of one division. But division can be avoided by applying the Legendre formula to count the number of occurrences of all prime factors from 2 to (p-1): 2. In both cases, the complexity of calculating all the coefficients from the sum (3) is equal to $\mathcal{O}(\frac{p-1}{2}\log_2^2 p)$.

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3: 4, 4	47: 44, 44	103: 100, 100	167: 164, 164
5: 8, 4	53: 40, 68	107: 108, 108	173: 200, 148
11: 12, 12	59: 60, 60	109: 104, 116	179: 180, 180
13: 8, 20	61: 72, 52	113: 124, 124	181: 200, 164
17: 12, 12	67: 68, 68	127: 124, 124	191: 188, 188
19: 20, 20	71: 68, 68	131: 132, 132	193: 204, 204
23: 20, 20	73: 76, 76	137: 156, 156	197: 200, 196
29: 40, 20	79: 76, 76	139: 140, 140	199: 196, 196
31: 28, 28	83: 84, 84	149: 136, 164	211: 212, 212
37: 40, 36	89: 76, 76	151: 148, 148	223: 220, 220
41: 28, 28	97: 76, 76	157: 136, 180	227: 228, 228
43: 44, 44	101: 104, 100	163: 164, 164	229: 200, 260
TABLE II			

THE TWISTED PAIRS OF EDWARDS CURVES

Squaring the calculated binomial coefficient $C_{\frac{p-1}{2}}^{j}$ also does not exceed $\mathcal{O}(\log_2^2 p)$.

Calculate all values of $d^j \mod p$ optimally applying recursive multiplication d^{j-1} on d. For this we make use of the Karatsuba multiplication method, which requires $\mathcal{O}(\log_2^{\log_2 3} p)$, rather than applying the Barrett method of modular multiplication. Therefore, the complexity of computing the entire tuple of degrees d^j , $j = 1, \ldots, n$ is $\mathcal{O}(\frac{p-1}{2}\log_2^{\log_2 3} p)$. Therefore, we obtain $\mathcal{O}(\frac{p-1}{2}\log_2^2 p)$.

Example: Number of curve points for d = 2 and p = 31 equals $N_{2[p]} = N_{2^{-1}[p]} = p - 3 = 28$.

Theorem 5: If $\left(\frac{d}{p}\right) = 1$, then the orders of the curves E_d and $E_{d^{-1}}$, satisfies to the following equality

$$|E_d| = |E_{d^{-1}}|.$$

If $\left(\frac{d}{p}\right) = -1$, then E_d and $E_{d^{-1}}$ are pair of twisted curves i.e. orders of curves E_d and $E_{d^{-1}}$ satisfies to the following relation of duality

$$|E_d| + |E_{d^{-1}}| = 2p + 2.$$

I.e. the orders of these curves complement each other up to 2p+2.

The Table II confirms this statement.

The first coordinate of the Table II element is p, which characteristic of the field F_p . Second and third coordinates are orders of curve E_d and correspond twist curve $E_{d^{-1}}$.

Proof: Let the curve be defined by $x^2 + y^2 = 1 + dx^2y^2 \pmod{p}$, then we can express y^2 in such way:

$$y^2 \equiv \frac{x^2 - 1}{dx^2 - 1} \pmod{p.}$$
 (22)

For $x^2 + y^2 = 1 + d^{-1}x^2y^2 \pmod{p}$ we could obtain that

$$y^2 \equiv \frac{x^2 - 1}{d^{-1}x^2 - 1} \pmod{p}$$
 (23)

Consider the case when $\left(\frac{d}{p}\right) = 1$ the case $\left(\frac{d}{p}\right) = -1$ proves analogously. if $\left(\frac{d}{p}\right) = 1$, then for the fixed x_0 a quantity of y over F_p can be calculated by the formula $\left(\frac{x^2-1}{d-1x^2-1}\right) + 1$ for x such that $d^{-1}x^2 + 1 \equiv 0 \pmod{p}$. For solution (x_0, y_0) to (7), we have

$$y^{2} = \frac{x^{2} - 1}{d^{-1}x^{2} - 1} = \frac{1 - x^{2}}{1 - d^{-1}x^{2}} = \frac{\left(\frac{1}{x^{2}} - 1\right)x^{2}}{\left(\left(\frac{d}{x^{2}}\right) - 1\right)d^{-1}x^{2}} = (24)$$

$$\left(\frac{1}{x^{2}} - 1\right)$$

$$=\frac{(\frac{1}{x^2}-1)}{((\frac{d}{x^2})-1)}d.$$

Thus, if (x_0, y_0) is solution of (7), then $\left(\frac{1}{x_0}, \frac{y_0}{\sqrt{d}}\right)$ is a solution to (23) because last transformations determines that $\frac{y_0^2}{d} \equiv \frac{d^{-1}\left(\frac{1}{x_0}\right)^2 - 1}{\left(\frac{1}{x_0}\right)^2 - 1} modp$. Therefore last transformations $(x_0, y_0) \rightarrow \left(\frac{1}{x_0}, \frac{y_0}{\sqrt{d}}\right) = (x, y)$ determines isomorphism and bijection.

Example: The number of points on E_d for d = 2 and $d^{-1} = 2$ with p = 31 is equal to $N_{2[31]} = N_{E_2^{-1}[31]} = p - 3 = 28$.

Example: If p = 7 and $d = 2^{-1} \equiv 4 \pmod{7}$, then we have $\left(\frac{d}{p}\right) = 1$ and the curve $E_{2^{-1}}$ has four points which are (0,1); (0,6); (1,0); (6,0), and the in case p = 7 for $d = 2 \pmod{7}$, the curve $E_{2^{-1}}$ also has four points which are (0,1); (0,6); (1,0); (6,0).

Proposition 2: If coefficients a, d of twisted Edwards curve $E_{a,d}$ satisfy the condition $\left(\frac{ad}{p}\right) = -1$ then this $E_{a,d}$ is isomorphic to Edwards curve $E_{d'}$. A twist curve to $E_{a,d}$ can be obtained due to multiplying on $c \in F_p$, $\left(\frac{c}{p}\right) = -1$.

In the case $\left(\frac{a}{p}\right) = 1$, $\left(\frac{d}{p}\right) = -1$ we obtain the isomorphism due to applying linear substitution $X = \frac{x}{\sqrt{a}}$, Y = y. This yields $x^2 + y^2 = 1 + \frac{d}{a}x^2y^2$ therefore we set $d' = \frac{d}{a}$ and obtain the final form $X^2 + Y^2 = 1 + d'X^2Y^2$. This curve have one point D = (0, -1) of order 2 and two points $\pm F = \left(\pm \frac{1}{\sqrt{a}}, 0\right)$ of order 4. These point has correspondent isomorphic images on Montgomery curve: $D_M = (0, 0)$ and $\pm F_M = (0, \pm \sqrt{a})$. Note that $\left(\frac{d'}{p}\right) = -1$. In the case $\left(\frac{a}{p}\right) = -1$, $\left(\frac{d}{p}\right) = 1$ we obtain the isomorphism due to new linear substitution $x \mapsto \frac{1}{X}$, $y \mapsto Y$. In result of this we have $\frac{1}{X^2} + Y^2 = 1 + d\frac{Y^2}{X^2}$ then, after the multiplying on X^2 , we finally get $1 + X^2Y^2 = X^2 + dY^2$.

To extend the result of previous Theorem 5: on twisted Edwards curve we consider the linear substitution $d \rightarrow cd$, $a \rightarrow ca$, where $\left(\frac{c}{p}\right) = -1$, transforms initial curve into twist curve.

Definition. We call the embedding degree a minimal power k of a finite field extension such that the group of points of the curve can be embedded in the multiplicative group of \mathbb{F}_{p^k} .

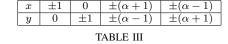
Let us obtain conditions of embedding [18] for the group of supersingular curves $E_d[\mathbb{F}_p]$ of order p in the multiplicative group of field \mathbb{F}_{p^k} whose embedding degree is k = 12 [18]. We now utilise the Zsigmondy theorem which implies that a suitable characteristic of field \mathbb{F}_p is an arbitrary prime p which do not divide 12 and satisfies the condition $q | \mathbf{P}_{12}(p)$, where $\mathbf{P}_{12}(x)$ is the cyclotomic polynomial. This p will satisfy the necessary conditions $(x^n - 1) \not| p$ for an arbitrary $n = 1, \ldots, 11$.

Proposition 3: The degree of embedding for the group of a supersingular curve E_d is equal to 2.

The proof is based on the fact that the order of the group of a supersingular curve E_d is equal to $p^k + 1$ and will be detailed in our next article.

We interesting the degree k needed for the Tate pairing embedding of the subgroup of order l (i.e. the smallest k such that $l | (q^k - 1)$ is relatively small.

Consider E_2 over F_{p^2} , for instance we assume p = 3. We define F_9 as $F_3(\alpha)$, where α is a root of $x^2 + 1 = 0$ over F_9 . Therefore elements of F_9 have form: $a + b\alpha$, where $a, b \in F_3$. So we assume that $x \in \{\pm(\alpha + 1), \pm(\alpha - 1), \pm\alpha\}$ and check its belonging to E_2 . For instance if $x = \pm(\alpha + 1)$ then $x^2 = \alpha^2 + 2\alpha + 1 = 2\alpha = -\alpha$. Also in this case $y^2 = \frac{2\alpha - 1}{\alpha - 1} = \frac{(2\alpha - 1)(\alpha + 1)}{(\alpha - 1)(\alpha + 1)} = \frac{(2\alpha - 1)(\alpha + 1)}{(\alpha - 1)(\alpha + 1)} = \frac{\alpha}{-2} = \alpha$. Therefore the correspondent second coordinate is $y = \pm(\alpha - 1)$. The similar computations lead us to full the following list of curves points.



POINTS OF EDWARDS CURVE OVER SQUARE EXTENSION.

The total amount is 12 affine points that confirms Corollary IV and Theorem IV because of $p^n - 3 - 2(-p)^{\frac{n}{2}} = 3^2 - 3 - 2(-3) = 12$.

5. Results

Our method for determining the order of a curve is valuable in itself and, in addition, provides a way to solve the inverse problem, that is, construct a curve of a given order. The conditions on a coefficients determining pairs of twisted curves of Edwards was found by us. We found a degree of embedding for the group of a supersingular curve in finite field with a minimal degree of extension. Due to Theorem 5. we obtain method to calculate the order of twisted curve for any Edwards curve. Let the first number denotes the quotient p (i.e. mod p). The second number is the number of points for d = 2 and the third number is the number solution for $d = 2^{-1} = (p+1)/2$.

According to GlobalSign, elliptical curve cryptography (ECC) can be used on most of today's modern browsers and operating systems. The list below shows which OS and browser versions are known to be compatible with ECC: Apple OS X, OS X 10.6, Google Android 4.0 Microsoft Windows and many other. One of the fundamental problems in EC cryptography is the generation of cryptographically secure ECs over prime fields, suitable for use in various cryptographic applications is solved due our method. A typical requirement of all such applications is that the order of the EC [5], [7], [18], [22]. One of essential requirement for EC is its order

(number of elements in the algebraic structure induced by the EC) possesses certain properties (e.g., robustness against known attacks [6, 16], small prime factors [5], [18], etc), which gives rise to the problem of how such EC can be generated. One of good decision of this tusk is a curve of big prime order [5], [18], [20]. Our method of determination of elliptic and Edwards curve order give one of possible decisions of these problems by analytic expression for the of curves order over a finite field F_{p^n} .

It is very important to avoid of curves of order p+1 because of it is tractable to the pairingbased attacks [5]. Our Theorems IV and IV establish criterion and sufficient conditions for such curves with order p + 1 therefore, it completely solves this problem. Additionally conditions of embedding of such curves was obtained. That is give rise to using the Edwards curve in the methods of zero knowledge proof such as Zk-Snark and Zk-Stark.

6. Discussions

The complexity of the discrete logarithm problem in the group of points of an elliptic curve depends on the order of this curve. Therefore, for intelligent control of the security level in the system, it is important to apply our method of determining the order [7], [8], [16], [21], of the curve providing the necessary level of cryptographic stability of the system.

One of most perspective problems arises on base of classes of supersigular curves studied by us is finding supersingular curves with small noninteger endomorphism [37]. Also, the further task is to investigate the stability of group of our curves against the Frey-Ruke attack [34]. Thus the appropriate task of study of these curves is to find the subgroups of divisors of prime order l and the quotient group $^{G}/_{lG}$, where $G = Pic_{E_d}^0(F_{q^k})$. Recently Boneh and Franklin developed identity-based cryptosystem [31] using the Weil pairing for the homomorphism in finite field. The next perspective problem will be the overcoming of constructing of supersingular curves of determined order with degree of embedding no greater than 6. Also the task that evoking interest and satisfying requirements of modern identity-based cryptosystems is to find the least embedding degree k for Montgomery curves with a given discriminant. Also one of the perspective goals is the finding of ordinary curves satisfying the following conditions of Tate paring:

1. There is a large prime l dividing the order of the group $G = Pic_{E_d}^0(F_{q^k})$.

2. The degree k needed for the Tate pairing embedding of the subgroup of order l (i.e. the smallest k such that $l \mid (q^k - 1)$ is relatively small.

Note that classes of supersingular curves, found by us, satisfy aforementioned requirements with extention degree 2. Number of Cryptographic block-chain protocols use groups satisfying the following requirements:

- to use supersingular curves, of which the most efficient ones appear to be in characteristic three.

- to use, so-called, MNT curves in large prime characteristic [32]. These are ordinary, as opposed to supersingular curves.

There are a number of advantages in using supersingular curves. For example they possess distortion maps which makes various protocols possible [33] or enable proofs to work [33], in addition supersingular curves have very efficient algorithms in characteristic three for computing the Tate pairing [36] and the underlying field arithmetic can be implemented relatively efficiently. Indeed, there are many advantages in using supersingular curves.

7. Conclusions

In this paper we have found a new effective algorithm for counting the elliptic and Edwards curves order over a finite field. In addition, the criterion for supersingularity of these curves has been was obtained as a result of this the using of groups of points of curves of genus 1 in cryptography which is intractable to MOV attacks. We have investigated the number of affine and projective points. It should be noted that our results provide a possibility to determine the number of such points in a constructive way, which provides a way for some intelligent control of the security level in the cryptosystem based on EC. The embedding degree for the Edwards and Montgomery curves have also been determined. The criterion of the supersingularity of Montgomery curve over finite field is found. Singular points of twisted Edwards curve are completely described.

We proved using existence of the birational isomorphism between twisted Edwards curve and elliptic curve in Weierstrass normal form that the result about order of curve over finite field can be applied to cubic in Weierstrass normal form.

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