

# Generation of Random Key Stream using Word Grid Puzzle for the Applications of Cryptography

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**Abstract:** The amount of digital data created and shared via internet has been increasing every day. The number of security attacks and threats has been increased due to poor selection of secret keywords and passwords. Cryptographic algorithms and security protocols are primarily rely on random keys to provide security services. Random numbers and key stream are playing major role in applying the security mechanisms in real time. In this paper, a novel method to generate random key stream using word grid puzzle is proposed. The proposed method is experimented and a minor analysis has been performed in the obtained result.

**Keywords:** Cryptography, security protocols, Password, Random key stream, Word grid puzzle

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## 1. Introduction

With the fast progression of digital data exchange in electronic way, Information Security is becoming much more important in storage and transmission of digital data [5]. Cryptography is the practice and study of techniques for providing secure communication over an insecure channel [2]. In network security, cryptography has a long history by providing a way to store sensitive information or transmit it across insecure networks (i.e. the Internet). The strength of any cryptosystem and security mechanisms are mainly depends on selection of unpredictable random keys. Diffusion and confusion are the two major building blocks for any cryptographic system. Confusion seeks to make the relationship between the statistics of the

ciphertext and the value of the encryption key as complex as possible [1, 6]. Symmetric key algorithms are well accepted in the modern communication network. The advantage of symmetric key cryptography is that the key management is very simple. In case of symmetric key method, the key should never be revealed to other users and should be kept secure. The key should be known to sender and the receiver only and no one else [4]. Depending on the nature of the randomness source, generators are classified in two categories as follows [6].

A. True random number generators (TRNG), where the source is a natural physical phenomenon and the properties of independence and unpredictability of the generated values are guaranteed by physical laws.

B. Pseudo random number generators (PRNG), where the source of randomness is a random initial value, called seed, which is expanded by means of a deterministic recursive formula, providing a modality for generating random sequences using only software methods.

In this paper, a novel method to generate random key stream with word grid puzzle is proposed.

The rest of the paper is organized as follows: Section II provides the literature survey, Section III briefs about proposed random key generator. Section IV presents the experimental results and analysis. The paper is concluded in Section V.

## 2. Literature Survey

Each and every encryption algorithm may be block cipher or stream cipher. Both types typically need a random key to achieve its purpose. However, stream cipher needs random key stream of length equals to the information to be encrypted. Cryptanalysis techniques like linear cryptanalysis, n-gram analysis, meet in the middle attack, brute force attack, man in the middle attack are usually performed to identify the secret key [2]. The efficiency of the ciphers are being depends on their throughput and memory requirement. Using of large key spaces with several numbers of rounds with multiple complex operations provide needed security[3].

Users rarely choose passwords that are both hard to guess and easy to remember. To determine how to help users choose good passwords, in [7], the authors performed a controlled trial of the effects of giving users different kinds of advice.

Some guidelines to choose a best password are [passwordgenerator.net]: use a password

that has at least 16 characters, use at least one number, one uppercase letter, one lowercase letter and one special symbol. Do not use postcodes, house numbers, phone numbers, birth dates, ID card numbers, and social security number in your passwords. Do not store your critical passwords in the cloud. Encrypt and backup your passwords to different locations.

To provide timely feedbacks to users, every Internet service now imposes a password strength meter (PSM) upon user registration or password change. It is a rare bit of good news in password research that well-designed PSMs do help improve the strength of user-chosen passwords. When choosing passwords for a new web service, most users (77.38%) simply retrieve one of their existing passwords from memory and then reuse (or slightly modify) it [8]

The randomness comes from atmospheric noise is better than the pseudo-random number algorithms typically used in computer programs. The best data security practice is not to let anyone but yourself generate your most important passwords [random.org]. A method for the human-assisted generation and application of pseudo-random keys for the purpose of encoding and decoding digital watermarks to and from a digitized data stream is proposed in [US5822432A].

In [9], the authors presented a multilayer image encryption and decryption using random puzzle based method to embed secret data into the cover image. This approach enhances the security and provides robust embedding. An 8×8 Sudoku puzzle which has a 64×64 reference matrix is utilized. The method gives a superior average capacity of 5 bits per pixel. The message is embedded at the sender using random permutation puzzle. At the Receiver end, the user has to rearrange the image thereby solving the puzzle.

In [10], a steganography method based on data embedding is introduced by using Sudoku solution matrix. In this scheme RGB value of Cover-image contains both secret information and key (Sudoku). RG component contains Secret information and B component contains key value. Same Sudoku solution matrix is used for embedding and extraction phase.

The word grid puzzle is  $n \times n$  matrix which contains words in a jumbled manner. In the proposed method, a novel method to generate random key stream with word grid puzzle is proposed.

### 3. Proposed Random Key Stream Generator

This section presents the overall working model of the proposed random stream generation process. Random key stream generation involves two requirements. First requirement is the word grid puzzle which is  $n \times n$  matrix and it consists of words in a jumbled manner. Second requirement is the words that are to be searched in the given word grid. Once the words to be searched in the word grid is found then the particular cells in the grid is alone made as 0 or -1. This indicates that the concerned cells are visited. Further, various kinds of 2-D array scanning are performed on the grid to generate different random key streams of various length.

The sequence of operation involved in generating random key stream using word grid puzzle is shown in Figure 1.

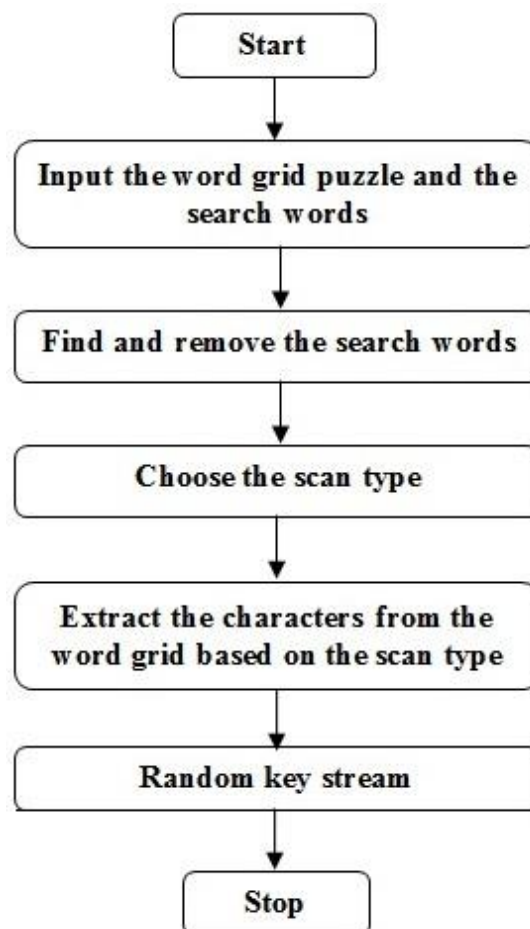


Figure. 1 Flowchart of proposed random key stream generation

### 3.1 Algorithm for Random Key Stream Generation

In this section, the sequence of steps involved in generation of random key stream is given.

**Input:** Word grid puzzle, Search words, Scan type      **Output:** Random key stream

Step 1: Start the process.

Step 2: Input the word grid puzzle and the corresponding search words.

Step 3: Find and remove the search words.

Step 4: Input the matrix scantype.

Step 5: Extract characters from the word grid based on the chosen scantype.

Step 6: Store the obtained characters as random key stream.

Step 7: Stop the process.

### 3.2 ScanTypes

Scan is the process of traversing all the cells of a 2-D matrix in a sequential pattern. The traversals can be an odd position scanning, even position scanning, snake scanning, zig-zag scanning, diagonal scanning, spiral scanning, etc. The sample scanning types are briefed below:

- **Odd scanning:** To generate random key streams by using the characters present in the odd location in the grid.
- **Even scanning:** To generate random key streams by using the characters present in the even location in the grid.
- **Diagonal scanning:** To generate random key streams by using the

characters present in the diagonal location in the grid.

- **Zig-Zag scanning:** To generate random key stream by using the characters present in the zigzag location in the grid.
- **Mid left:** To generate random key streams by using the characters present in the middle and towards the beginning of the grid. Similarly, Mid right.
- **Upper triangular scanning:** To generate random key streams by using the characters present in the upper triangular of the grid. Similarly, Lower triangular scanning.

The size of key stream that are to be generated from the grid depends on the size of the key mentioned in the algorithm. For instance, if the chosen algorithm is Data Encryption Standard, the 64 bits will be derived. Also, the kind of scanning and volume of the key stream are depends on the understanding between the sender and the receiver.

### 3.3 Illustration for the Proposed Method

In this section, the proposed random key stream generation method is illustrated with a sample word grid puzzle. The sample word grid is shown in Table 1. The search words of the word grid are Intellect, Infinity formula, Elite, Paratrooper and Integrity.

**Table 1. Sample Word Grid**

Q	B	D	H	I	Y	X	A	P	O	L	M	I	F	T	Y
D	A	P	U	L	M	K	G	F	E	F	S	N	H	D	B
S	C	N	Q	P	T	G	A	G	Y	O	A	F	U	T	G
G	N	F	R	O	G	S	G	T	Q	A	Q	I	I	C	Z
B	M	J	H	L	C	A	I	P	G	L	B	N	P	E	C
M	T	K	J	M	N	R	X	S	X	D	N	I	I	L	D
D	F	W	U	Z	G	D	B	O	J	E	Y	T	N	L	J
B	L	T	P	E	N	S	F	K	X	L	R	Y	R	E	Q
Y	J	U	T	Z	M	R	R	J	Y	I	U	F	F	T	P
U	B	N	D	B	I	B	S	H	M	T	O	O	U	N	L
P	I	I	B	J	T	Y	H	S	V	E	L	R	D	I	V
O	X	V	M	F	E	X	X	E	L	U	P	M	G	U	X
P	A	R	A	T	R	O	O	P	E	R	A	U	Q	Y	R
W	Z	Z	C	Q	H	Q	H	Y	D	U	Q	L	O	I	W
F	R	V	V	G	P	Z	O	M	H	M	I	A	D	H	U
A	T	B	N	O	P	R	H	F	M	Z	N	D	T	J	A

**Table 2. Word Grid after Highlighting the Search Words**

Q	B	D	H	I	Y	X	A	P	O	L	M	I	F	T	Y
D	A	P	U	L	M	K	G	F	E	F	S	N	H	D	B
S	C	N	Q	P	T	G	A	G	Y	O	A	F	U	T	G
G	N	F	R	O	G	S	G	T	Q	A	Q	I	I	C	Z
B	M	J	H	L	C	A	I	P	G	L	B	N	P	E	C
M	T	K	J	M	N	R	X	S	X	D	N	I	I	L	D
D	F	W	U	Z	G	D	B	O	J	E	Y	T	N	L	J
B	L	T	P	E	N	S	F	K	X	L	R	Y	R	E	Q
Y	J	U	T	Z	M	R	R	J	Y	I	U	F	F	T	P
U	B	N	D	B	I	B	S	H	M	T	O	O	U	N	L
P	I	I	B	J	T	Y	H	S	V	E	L	R	D	I	V
O	X	V	M	F	E	X	X	E	L	U	P	M	G	U	X
P	A	R	A	T	R	O	O	P	E	R	A	U	Q	Y	R
W	Z	Z	C	Q	H	Q	H	Y	D	U	Q	L	O	I	W
F	R	V	V	G	P	Z	O	M	H	M	I	A	D	H	U
A	T	B	N	O	P	R	H	F	M	Z	N	D	T	J	A

Further, find the search words in the word grid and the word grid after removing the search words is shown in Table 2. The

remaining characters present in the word grid (matrix) are utilized to generate random key streams. Also, different types of

scanning are can be performed in the grid to generate random key streams of any required size. Each scanning paves way for a unique key stream for the same grid.

The random key streams obtained by using few of the scan types are given below:

▪ **Key stream obtained using odd scan:**

QDIXPLITDPLKFFNDSNPMGOF  
TGFOSTAICBJLAPLNEMKMRS  
ILDWZDOETLBTESKLYEYUZRJ  
IFTUNBBHTONPIJYSERIOVFXE  
UMUPRTOPRUYWZQQYULIFVG  
ZMMAHABORFZDJ

▪ **Key stream obtained using Zig-Zag**

scan:QBDHIYXAPOLMIFTYBDH  
NSFEFGKMLUPASCNQPTMAGY  
OAFUTGZCIIQAQTGSGORFNB  
MJHLCAIPGLBNPECDLIINDXSX  
RNMJKTDFWUZGDBOJEYTNLJ  
QERYRLXKFSNEPTLYJUTZMRR  
JYIUFFTPLNUOOTMHSBIBDNB  
PIIBJTYHSVELRDIVXUGMPULE  
XXEFMVXPARATROOPERAUQ  
YRWIOLQUDYHQHCZZFRVV  
GPZOMHMIADHUAJTDNZMFHR  
PONBT

▪ **Key stream obtained using spiral**

scan:QBDHIYXAPOLMIFTYBGZ  
CDJQPLVXRWUAJTDNZMFHRP  
ONBTAFWPOPUIYBDMBGSDAP  
ULMKGFEFSNHDTCLELETNIUY  
IHDAIMHMOZPGVVRZAXIBJLF  
TMNCNQPTMAGYOAFUIPINRF  
UDGQOLQUDYHQHCZRVINU  
TWKJFROSGTQAQINITYFORM  
UAREPOORTAMBDTPUJHLCAIP  
GLBNYRUOLPULEXXEFJBZEZ  
MNRXSXDELITEVSHYTIMNGD  
BOJXYMHSBRSFKJR

Similar to this, various other types of scanning can be applied on the grid to generatedifferent and unique random key streams. Scan which produces less amount of bit streams can be utilized to generate keys for play fair, DES, and AES algorithms. Scan which produces more amount of bit streams can be utilized to generate keys for stream ciphers.

### 3.4 Salient Features of the Proposed Method

The following are the salient features of the proposed method:

- The proposed method can be used to generate keys of variable size. Hence, it can be utilized to generate keys for both block and stream ciphers.
- The Word Grid puzzle and the words present in the grid are known to the communicating persons. It may not give any suspicion to the intruders/attackers about the purpose of sharing the puzzles.
- The overhead associated with key generation, key distribution and keeping the key secure can be reduced.
- The generated key stream can be used for padding, initialization vectors, and salt values, session keys, etc.

## 4. Experimental Results And Analysis

The proposed method is experimented using Python language and the system configuration is Processor Intel Core i3 CPU, Clock speed 3.07 GHz, RAM 4GB and the operating system is Windows (64bit). The experimental result of the proposed method is given in Table 3 for few scan types. From the result, it is observed that the bit stream is

random and hence suitable to generate keys for cryptographic applications.

**Table 3. Experimental results of the proposed method**

Scan	Random key stream in binary format
Odd	101000110001001001001101100010100001001100100100110101001000100101000010011001 001011100011010001101001110100010010100111001110101000010011011000111100111110 001101010100100111110100111010100100000110010011000011100001010010101001100100 000110100001001100100111010001011001101100101110011011010010101001110001001001 001100110010001001010111101101010001001001111100010110101001001100100001010101 001000101101001110010111001100101100110001011011001101010110110101010010100101 01001001100011010101001010101001110100001010000101001000101010010011111001110 10100001001001100101010100110100111000101101001010010011001111101011010001101 0110001000101101010110011011010101101000010100101010010011111010000101001010 101011011001101011110110101010001101000110110011010101100110010010011000110101 01101000111101101010011011001101000001100100010000011000010100111110100101000 110101101010001001001010
Zig-Zag	101000110000101000100100100010010011011001101100010000011010000100111110011001 001101100100110001101010100111100110000101000100100100010011101010011100011010 001011000110100011110010111001101100110010101011010000100000110100111000011100 11101010001101000010101001001101100000110011110110011001111100000110001101010 101101010010001111011010100001110010011001001101000110000011010001101010010001 111010011100011110011111010010100011010011101000010100110110010101001000100110 010000111000001100100110100001000111100110010000101001110101000010001011000011 100010010011001001001100100110011101000100101100010100111011000101001010011101 00110110010101001011101010010001001000110101011101010110110101000111100010010 000101001111100101010001011011001101010010011101001100100101010100011000101101 001010110011010010100110010110001001011100011010100111001110100010110100001010 100100110010110011001010101010110101001011010100110110100101010010100101010110 011001001101010110001101000110101010010100001001100100111010101011001111100111 110101001001101100100010100111000010100100110000101000100100111010000101010000 100100110010011000010100101010101001011001100100010100111010110100010110011001 010010100010010010011010110101100010101011000111100110110100001010101100110010 001011011000101100010001011000110100110110101101011000101000010000011010010100 000110101001010010100111110011111010000100010110100101000001101010110100011011 00110100101010111001001100111110011001010001101010110001001011001100100010100 0110010001010001100001110110101011010100011010100101011010101101000111101000 010110101001111100110110010001001101100100110000011000100100100010101011000001 100101010101001000100100111010110101001101100011010010001010010101000010011111 00111010000101010100
Upper triangle	101000110000101000100100100010010011011001101100010000011010000100111110011001 001101100100110001101010100111100110000011010000101010110011001001101100101110 001111000110100010110001101010011100111010010001000100100001010011101010001101 000010101001001101100000110001111011001100111110000011000110101010110101001000 111101001010011111000111101001110001111010100101000110000011010001100100110010 011000011101101010011001000011100000110010011010000100011110011001000010100111 010100001000101100001110011101010010101100010100111011000100010010011101001001 100100110011001000100100010010000101001111100101010001011011001101010010011101 001100100101010001101001011101100010011001010010101100110100101000101101000110 010101011001100100110101011000110100011010101001010000100110110101001001111100 111110101011001110100110010001011001100101001010001001001001101011010100001001 1011000111101010110110001010101010001101100110100101001111100100110101110010 001010101100001

The execution time taken to generate key streams by using various scan types is given in Table 4. From the table, it is inferred that the execution time to generate the bit

streams is vary less and hence it is suitable for real time applications.

**Table 4. Execution time**

S.No.	Scan type	Time take (seconds)
1.	Odd	0.4266
2.	Even	0.4489
3.	Zig-Zag	0.8353
4.	Spiral	0.8433
5.	Upper triangle	0.8746

### 5. Conclusion and Future Work

In this paper, a novel and simple method to generate the random key stream using Word Grid Puzzle is developed. The proposed method is a new notion to generate random key streams. The generated bit stream is random and the execution time is less than 1 seconds. This new methodology for generating key stream using Word Grid puzzle is an effective and secure method by confusing the attackers about the actual purpose of sharing the puzzles. In future, randomness testing is performed with NIST SP 800-22 test cases to access the randomness of the proposed random number generator. To apply the generated random numbers to encrypt/decrypt images.

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